# Matrix compositions

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# Motivation: L-convex polyominoes

A polyomino is a connected finite set of cells in the plane  $\mathbb{Z} \times \mathbb{Z}$ . Some examples are



A polyomino is *L*-convex when every pair of cells can be connected by an internal path with at most one change of direction. Using formal power series techniques, in

G. Castiglione, A. Frosini, A. Restivo, S. Rinaldi, *Enumeration of L-convex polyominoes*, 2005

it was proved that the numbers  $f_n$  of all Lconvex polyominoes with perimeter 2(n + 2)satisfy the recurrence

 $f_{n+2} = 4f_{n+1} - 2f_n \qquad (n \ge 1).$ Specifically  $f_n = 1, 2, 7, 24, 42, 120, \dots$  Compositions and 2-compositions

A composition of a number  $n \in \mathbb{N}$  is any k-tuple  $(x_1, \ldots, x_k)$  of positive integers such that

 $x_1 + \dots + x_k = n.$ 

A 2-composition of length k of a number  $n \in \mathbb{N}$  is a  $2 \times k$  matrix

$$M = \begin{bmatrix} x_1 & x_2 & \dots & x_k \\ y_1 & y_2 & \dots & y_k \end{bmatrix}$$

with nonnegative integer entries, without zero columns, such that the sum of all the entries is n.

For instance the 2-compositions of n = 2 are:

$$\begin{bmatrix} 0\\2 \end{bmatrix}, \begin{bmatrix} 1\\1 \end{bmatrix}, \begin{bmatrix} 2\\0 \end{bmatrix}, \begin{bmatrix} 1&0\\0&1 \end{bmatrix}, \begin{bmatrix} 0&1\\1&0 \end{bmatrix}, \begin{bmatrix} 1&1\\0&0 \end{bmatrix}, \begin{bmatrix} 0&0\\1&1 \end{bmatrix}.$$

The number of all 2-compositions of n is  $f_n$ 

## Polyominoes and 2-compositions

Since the number of all *L*-convex polyominoes with perimeter 2(n+2) is equal to the number all 2-compositions of *n*, there exists a bijection between this two combinatorial classes.

In

G. Castiglione, A. Frosini, E. Munarini, A. Restivo, S. Rinaldi, *Enumeration* of *L*-convex polyominoes. *II. Bijection* and area, FPSAC'05

it was given an explicit bijection.

Here follows a sketch.

## The bijection



The factorization in  $c^h a$ ,  $bd^k$   $(h, k \ge 0)$ ,  $c^r d^s$   $(r, s \ge 1)$  implies the bijection:

$$c^{h}a \to \begin{bmatrix} h+1\\0 \end{bmatrix}, \ bd^{k} \to \begin{bmatrix} 0\\k+1 \end{bmatrix}, \ c^{r}d^{s} \to \begin{bmatrix} r\\s \end{bmatrix}$$
$$a(c^{0}a)(bd^{0})(cd)(c^{0}a)(bd^{0})(cd)(bd)(c^{0}a)b \rightsquigarrow$$
$$\sim \begin{bmatrix} 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1\\ 0 & 1 & 1 & 0 & 1 & 1 & 2 & 0 \end{bmatrix}.$$

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Aim of the talk: to give an elementary introduction to matrix compositions.

Structure of the talk:

- definition of matrix compositions
- ◊ recurrences and explicit form

 Cassini-like identity and combinatorial interpretation

- ◊ some encoding
- $\diamond$  other results.

### Definition of *m*-compositions

Let  $m \in \mathbb{N}$ , m > 0. An m-row matrix composition (m-composition for short) of *length* k of a number  $n \in \mathbb{N}$  is an  $m \times k$  matrix

$$M = \begin{bmatrix} x_{11} & \dots & x_{1k} \\ \vdots & & \vdots \\ x_{m1} & \dots & x_{mk} \end{bmatrix}$$

such that

$$\diamond \quad x_{ij} \in \mathbb{N} \qquad \forall i, j$$

$$\diamond \quad (x_{1j}, \dots, x_{kj}) \neq (0, \dots, 0) \quad \forall j$$

$$\diamond \quad \sigma(M) = \sum_{i=1}^{m} \sum_{j=1}^{k} x_{ij} = n$$

$$\mathcal{C}_n^{(m)} = \{ m \text{-compositions of } n \}$$

 $\mathcal{C}_{n,k}^{(m)} = \{ \text{ } m \text{-compositions of } n \text{ of length } k \}$ 

$$c_n^{(m)} = |\mathcal{C}_n^{(m)}| \qquad c_{n,k}^{(m)} = |\mathcal{C}_{n,k}^{(m)}|$$

#### **Multisets**

A multiset on a set X is a function  $\mu : X \to \mathbb{N}$ The multiplicity of an element  $x \in X$  is  $\mu(x)$ The order of  $\mu$  is  $\operatorname{ord}(\mu) = \sum_{x \in X} \mu(x)$ The number of all multisets of order k on a set of size n is the multiset coefficient

$$\binom{n}{k} = \frac{n^{\overline{k}}}{k!} = \frac{n(n+1)\dots(n+k-1)}{k!}$$

The columns of an m-composition are multisets with positive order on an m-set.

$$M = \begin{bmatrix} \vdots & c_1 & \vdots \\ \vdots & \vdots & \vdots \\ \vdots & c_m & \vdots \end{bmatrix}$$

Each column  $(c_1, \ldots, c_k)$  is equivalent to the multiset  $\mu : [m] \to \mathbb{N}$  defined by

$$\mu = \begin{pmatrix} 1 & 2 & \dots & m \\ c_1 & c_2 & \dots & c_k \end{pmatrix}$$

with positive order:  $\mu(1) + \ldots + \mu(m) > 0$ .

## Sum recursions

$$M = \begin{bmatrix} x_{11} & x_{12} & \dots & x_{1k+1} \\ \vdots & \vdots & & \vdots \\ x_{m1} & x_{m2} & \dots & x_{mk+1} \end{bmatrix}$$

$$c_{n+m,k+1}^{(m)} = \sum_{i=1}^{n+m-k} \binom{m}{i} c_{n+m-i,k}^{(m)}$$

$$c_{n+m}^{(m)} = \sum_{i=1}^{n+m} \binom{m}{i} c_{n+m-i}^{(m)}$$

#### Main recurrence

Let  $A_i$  be the set of all *m*-compositions

$$M = \begin{bmatrix} \vdots & \dots \\ x_{i1} & \dots \\ \vdots & \dots \end{bmatrix}$$

of n + m with  $x_{i1} \neq 0$ . Then

$$\mathcal{C}_{n+m}^{(m)} = A_1 \cup \ldots \cup A_m$$

and for the Principle of Inclusion-Exclusion

$$c_{n+m}^{(m)} = \sum_{\substack{S \subseteq [m] \\ S \neq \emptyset}} (-1)^{|S|-1} \left| \bigcap_{i \in S} A_i \right|.$$

 $\bigcap_{i\in S} A_i$  is formed by all the m-compositions

$$M = \begin{bmatrix} \vdots & \dots & \\ x_{i1} & \dots & \\ \vdots & \dots & \end{bmatrix} \sim \begin{bmatrix} \vdots & \dots & \\ x_{i1} - 1 & \dots & \\ \vdots & \dots & \end{bmatrix}$$

for every  $i \in S$ . Then

$$\left|\bigcap_{i\in S} A_i\right| = 2c_{n+m-|S|}^{(m)}.$$

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Therefore we have the main recurrence

$$c_{n+m}^{(m)} = 2\sum_{i=1}^{m} {m \choose i} (-1)^{i-1} c_{n+m-i}^{(m)}$$

For m = 2, 3, 4 we have the recurrences:

$$c_{n+2}^{(2)} = 4c_{n+1}^{(2)} - 2c_n^{(2)}$$
$$c_{n+3}^{(3)} = 6c_{n+2}^{(3)} - 6c_{n+1}^{(3)} + 2c_n^{(3)}$$
$$c_{n+4}^{(4)} = 8c_{n+3}^{(4)} - 12c_{n+2}^{(4)} + 8c_{n+1}^{(4)} - 2c_n^{(4)}.$$

Similarly

$$c_{n+m,k+1}^{(m)} = \sum_{i=1}^{m} {m \choose i} (-1)^{i-1} c_{n+m-i,k}^{(m)} + \sum_{i=1}^{m} {m \choose i} (-1)^{i-1} c_{n+m-i,k+1}^{(m)}$$

#### Explicit form

Let  $A_i$  be the set of all matrices  $M \in \mathcal{M}_{m,k}(\mathbb{N})$ where the *i*-th column is the zero vector and  $\sigma(M) = n$ . Then

$$\mathcal{C}_{n,k}^{(m)} = A_1' \cap \ldots \cap A_k'$$

and for the Principle of Inclusion-Exclusion

$$c_{n,k}^{(m)} = |A'_1 \cap \ldots \cap A'_k| = \sum_{S \subseteq [k]} (-1)^{|S|} \left| \bigcap_{i \in S} A_i \right|.$$

An element in  $\bigcap_{i \in S} A_i$  is a matrix  $M \in \mathcal{M}_{m,k}(\mathbb{N})$ with a zero vector in each column indexed by the elements of S. Removing such columns we just have a multiset of order n on a set of size mk - m|S|. Then

$$\bigcap_{i\in S} A_i = \left( \binom{m(k-|S|)}{n} \right)$$

and consequently

$$c_{n,k}^{(m)} = \sum_{i=0}^{k} \binom{k}{i} \binom{m(k-i)}{n} (-1)^{i}$$

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### Cassini-like identities

The numbers  $f_n = c_n^{(2)}$  of all 2-compositions of n satisfy a Cassini-like identity:

$$f_n f_{n+2} - f_{n+1}^2 = -2^{n-1}$$

for every  $n \geq 1$ .

Equivalently

$$\begin{vmatrix} f_n & f_{n+1} \\ f_{n+1} & f_{n+2} \end{vmatrix} = -2^{n-1}.$$

For the numbers  $c_n^{(m)}$  such an identity becomes

$$\begin{vmatrix} c_n^{(m)} & c_{n+1}^{(m)} & \dots & c_{n+m-1}^{(m)} \\ c_{n+1}^{(m)} & c_{n+2}^{(m)} & \dots & c_{n+m}^{(m)} \\ \vdots & \vdots & & \vdots \\ c_{n+m-1}^{(m)} & c_{n+m}^{(m)} & \dots & c_{n+2m-2}^{(m)} \end{vmatrix} = (-1)^{\lfloor \frac{m}{2} \rfloor} 2^{n-1}$$

for every  $m \geq 1$  and  $n \geq 1$ .

#### First reduction

Let

$$C_{n}^{(m)} = \begin{bmatrix} c_{n}^{(m)} & c_{n+1}^{(m)} & \dots & c_{n+m-1}^{(m)} \\ c_{n+1}^{(m)} & c_{n+2}^{(m)} & \dots & c_{n+m}^{(m)} \\ \vdots & \vdots & & \vdots \\ c_{n+m-2}^{(m)} & c_{n+m-1}^{(m)} & \dots & c_{n+2m-3}^{(m)} \\ c_{n+m-1}^{(m)} & c_{n+m}^{(m)} & \dots & c_{n+2m-2}^{(m)} \end{bmatrix}$$

By the main recurrence

$$c_{n+m}^{(m)} = \alpha_{m-1}c_{n+m-1}^{(m)} + \dots + \alpha_1c_{n+1}^{(m)} + \alpha_0c_n^{(m)}$$

where

$$\alpha_k = (-1)^{m-k-1} 2\binom{m}{k}.$$

The last row can be simplified subtracting a suitable linear combination of the first m-1 rows.

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We obtain the determinant



Extracting  $\alpha_0 = (-1)^{m-1}2$  from the last row and shifting cyclically all rows downward we have

$$\det C_n^{(m)} = 2 \det C_{n-1}^{(m)}.$$

Then, for every  $n \ge 1$ , it follows that:

$$\det C_n^{(m)} = 2^{n-1} \det C_1^{(m)}$$

where

$$C_1^{(m)} = [c_{i+j+1}^{(m)}]_{i,j=0}^{m-1}.$$

To compute the determinant of this matrix it is useful to consider a combinatorial interpretation of m-compositions.

Combinatorial interpretation

Given the 3-composition

$$M = \begin{bmatrix} 2 & 0 & 1 & 2 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 2 \end{bmatrix}$$

consider its entries as the number of occurrences of three given colors, for instance red, green and blue, linearly ordered red < green < blue. Then

$$M = \begin{bmatrix} 2 & 0 & 1 & 2 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 2 \end{bmatrix}$$

is equivalent to a  $\underline{3-colored \ bargraph}$  of area 11 with 4 columns



Now, writing the columns horizontally, from the bottom to the top, we have a 3-colored linear partition of the set  $\{1, 2, ..., 11\}$  with 4 blocks:



In general, the m-compositions of n of length k are equivalent to the

- $\diamond$  *m*-colored bargraphs of area *n* with *k* columns.
- $\diamond$  *m*-colored linear partitions of [*n*] with *k* blocks;

#### A useful sum

Let  $\pi$  be an *m*-colored linear partition of

$$L = \{x_1, \dots, x_{i+1}, \dots, x_{i+j+1}\}.$$

The element  $x_{i+1}$  belongs to a block of the form  $\{x_{i-h+1}, \ldots, x_i, x_{i+1}, x_{i+2}, \ldots, x_{i+k+2}\}$ :

Removing this block,  $\pi$  splits into two *m*-colored linear partitions  $\pi_1$  and  $\pi_2$ . Then it follows that

$$c_{i+j+1}^{(m)} = \sum_{h,k\geq 0} \binom{m}{h+k+1} c_{i-h}^{(m)} c_{j-k}^{(m)}$$

#### Second reduction

The previous identity implies the decomposition

$$C_1^{(m)} = L^{(m)} M^{(m)} L_T^{(m)}$$

where

$$L^{(m)} = \begin{bmatrix} c_{i-j}^{(m)} \end{bmatrix}_{i,j=0}^{m-1} = \begin{bmatrix} c_0 & & & \\ c_1 & c_0 & & \\ \vdots & \ddots & \ddots & \\ c_{m-1} & \dots & c_1 & c_0 \end{bmatrix}$$
$$M^{(m)} = \begin{bmatrix} \begin{pmatrix} m & \\ i+j+1 \end{pmatrix} \end{bmatrix}_{i,j=0}^{m-1}.$$

Then

$$\det C_1^{(m)} = \det M^{(m)}.$$

To calculate det  $M^{(m)}$  we will use another identity coming form our combinatorial interpretation.

#### Another useful sum

Recall that  $\binom{m}{i+j+1}$  gives the number of all the order maps  $f: [i+j+1] \rightarrow [m]$ .

Suppose that f(i + 1) = k, with  $k \in [m]$ . Since f is order preserving, it follows that

 $\diamond \quad 1 \leq x \leq i \text{ then } 1 \leq f(x) \leq k$ 

 $\diamond \quad i+2 \leq x \leq i+j+1 \text{ then } k \leq f(x) \leq m.$ 

Then

$$\begin{pmatrix} m \\ i+j+1 \end{pmatrix} = \sum_{k=1}^{m} \binom{k}{i} \binom{m-k+1}{j}$$
$$= \sum_{k=0}^{m-1} \binom{i+k}{i} \binom{m-k}{j}$$

#### Final step

The previous identity implies that  $M^{(m)} = B^{(m)} \widetilde{B}^{(m)}$ 

where

$$B^{(m)} = \left[ \binom{i+j}{i} \right]_{i,j=0}^{m-1}$$

and

$$\widetilde{B}^{(m)} = \left[ \begin{pmatrix} m-i \\ j \end{pmatrix} \right]_{i,j=0}^{m-1}.$$

Since  $\tilde{B}^{(m)} = J^{(m)}B^{(m)}$  where

$$J^{(m)} = [\delta_{i+j,m-1}]_{i,j=0}^{m-1} = \begin{bmatrix} & & 1 \\ & & 1 \\ 1 & & \\ 1 & & \end{bmatrix},$$

we have

$$M^{(m)} = B^{(m)}J^{(m)}B^{(m)}$$

and

$$\det M^{(m)} = \det J^{(m)} (\det B^{(m)})^2.$$

Since, as well known,

$$\det J^{(m)} = (-1)^{\lfloor m/2 \rfloor}$$

and

$$\det B^{(m)} = 1$$

we have

$$\det M^{(m)} = (-1)^{\lfloor m/2 \rfloor}$$

and consequently

$$\det C_1^{(m)} = \det M^{(m)} = (-1)^{\lfloor m/2 \rfloor}.$$
  
Finally, since for every  $n \ge 1$  we have

$$\det C_n^{(m)} = 2^{n-1} \det C_1^{(m)},$$

then

$$\det C_n^{(m)} = (-1)^{\lfloor m/2 \rfloor} 2^{n-1}.$$

#### *m*-compositions as words

Consider each m-composition as the concatenation of its columns:

$$M = \begin{bmatrix} 1 & 3 & 2 & 1 \\ 0 & 2 & 1 & 0 \\ 1 & 0 & 2 & 3 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} \begin{bmatrix} 3 \\ 2 \\ 0 \end{bmatrix} \begin{bmatrix} 2 \\ 1 \\ 2 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 3 \end{bmatrix}$$

Each column  $(c_1, \ldots, c_k)$  is equivalent to the multiset  $\mu : [m] \to \mathbb{N}$  defined by

$$\mu = \begin{pmatrix} 1 & 2 & \dots & m \\ c_1 & c_2 & \dots & c_k \end{pmatrix}$$

with positive order:  $\mu(1) + \ldots + \mu(m) > 0$ .

This means that

$$\mathcal{C}^{(m)} = \{a_{\mu} : \mu \in \mathcal{M}_{\neq 0}^{(m)}\}^*$$

where  $\mathcal{M}_{\neq 0}^{(m)}$  is the set of all multisets  $\mu$ :  $[m] \to \mathbb{N}$  with positive order.

Let  $X = \{x_{\mu} : \mu \in \mathcal{M}_{\neq 0}^{(m)}\}$ . Then, for  $a_{\mu} \rightsquigarrow x_{\mu}$ , we have the generating series for  $\mathcal{C}^{(m)}$ 

$$c(X) = \frac{1}{1 - \sum_{\mu \in \mathcal{M}_{\neq 0}^{(m)}} x_{\mu}}$$

In particular, for  $x_{\mu} = x^{\operatorname{ord}(\mu)}$  we get

$$c^{(m)}(x) = \sum_{n \ge 0} c_n^{(m)} x^n = \frac{1}{1 - h(x)}$$

where

$$h(x) = \sum_{\mu \in \mathcal{M}_{\neq 0}^{(m)}} x_{\mu} = \sum_{k \ge 1} {\binom{m}{k}} x^{k}$$

that is

$$h(x) = \frac{1}{(1-x)^m} - 1.$$

Then

$$c^{(m)}(x) = \frac{(1-x)^m}{2(1-x)^m - 1}$$

From here we have the previous results and a the following new recurrence

$$c_{n+1}^{(m)} = -\delta_{n,0} + 2c_n^{(m)} + \sum_{k=0}^n {m+k-1 \choose k+1} c_{n-k}^{(m)}$$

### *m*-compositions of Carlitz type

A Carlitz composition is a composition without two equal consecutive elements.

We say that an m-composition is of Carlitz type when no two adjacent columns are equal.

For instance

$$M = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 2 & 2 & 4 & 2 \\ 4 & x & 3 & 3 \end{bmatrix}$$

is of Carlitz type only when  $x \neq 4$ .

Let  $\mathcal{Z}$  be the set of all words corresponding to the *m*-compositions of Carlitz type and let  $\mathcal{Z}_{\mu}$ be the subset of  $\mathcal{Z}$  formed by the words ending with  $a_{\mu}$ , for every  $\mu \in \mathcal{M}_{\neq 0}^{(m)}$ . Then

$$\mathcal{Z} = 1 + \sum_{\mu \in \mathcal{M}_{\neq 0}^{(m)}} \mathcal{Z}_{\mu}$$

 $\mathcal{Z}_{\mu} = (\mathcal{Z} - \mathcal{Z}_{\mu})a_{\mu} \qquad \forall \ \mu \in \mathcal{M}_{\neq 0}^{(m)}.$ 

To obtain the generating series for  $\mathcal{Z}$  and  $\mathcal{Z}_{\mu}$  substitute each letter  $a_{\mu}$  with the indeterminate  $x_{\mu}$ :

$$z(X) = 1 + \sum_{\mu \in \mathcal{M}_{\neq 0}^{(m)}} z_{\mu}(X)$$

$$z_{\mu}(X) = (z(X) - z_{\mu}(X))x_{\mu} \qquad \forall \ \mu \in \mathcal{M}_{\neq 0}^{(m)}.$$

Then

$$z_{\mu}(X) = \frac{x_{\mu}}{1 + x_{\mu}} z(X)$$

and finally

$$z(X) = \frac{1}{1 - \sum_{\substack{\mu \in \mathcal{M}_{\neq 0}^{(m)}}} \frac{x_{\mu}}{1 + x_{\mu}}}.$$

For  $x_{\mu} = x^{\operatorname{ord}(\mu)}$  we have

$$\sum_{n \ge 0} z_n^{(m)} x^n = \frac{1}{1 - \sum_{k \ge 1} \binom{m}{k} \frac{x^k}{1 + x^k}}$$

where  $z_n^{(m)}$  is the number of all *m*-compositions of Carlitz type of *n*.

Expanding this series we can obtain:

$$z_n^{(m)} = \sum_{k \ge 0} \sum_{\substack{\alpha, \beta \in \mathbb{N}_0^k \\ \alpha \cdot \beta = n}} {\binom{m}{\alpha}} ((-1)^{|\beta| - k}.$$

where

$$\begin{aligned} \alpha \cdot \beta &= a_1 b_1 + \dots + a_k b_k \\ |\beta| &= b_1 + \dots + b_k \\ \begin{pmatrix} m \\ \alpha \end{pmatrix} &= \begin{pmatrix} m \\ a_1 \end{pmatrix} \dots \begin{pmatrix} m \\ a_k \end{pmatrix} \end{aligned}$$
for every  $\alpha = (a_1, \dots, a_k)$  and  $\beta = (b_1, \dots, b_k).$ 

A regular language for *m*-compositions

The encoding used in

# A. Björner, R. Stanley, *An analogue of Young's lattice for compositions*, FP-SAC'05

for ordinary compositions can be extended to m-compositions as follows. For instance

More precisely, given the finite alphabet

 $\mathcal{A}_m = \{a_1, \cdots, a_m, b_1, \ldots, b_m\},\$ 

we can define a map  $\ell: \mathcal{C}^{(m)} \to \mathcal{A}_m^*$  setting



and proceeding as in the example.

The words of the language  $\mathcal{L}_m = \ell(\mathcal{C}^{(m)}) \subseteq \mathcal{A}_m^*$ , corresponding to the *m*-compositions, are characterized by the following conditions:

i) the first letter is  $a_1$  or ... or  $a_m$ ;

ii) each letter  $a_i$  or  $b_i$  can be followed by any  $b_j$ , while it can be followed by a letter  $a_j$  only when  $i \leq j$ .

This characterization implies a unique factorization of the form xy, where:

Then  $\mathcal{L}_m$  is a regular language defined by the unambiguous regular expression:

 $\varepsilon + \mathcal{L}_m' \mathcal{L}_m''$ 

where  $\varepsilon$  is the empty word and

$$\mathcal{L}'_m = a_1^+ a_2^* \dots a_m^* + a_2^+ a_3^* \dots a_m^* + \dots + a_m^+,$$
$$\mathcal{L}''_m = (b_1 a_1^* a_2^* \dots a_m^* + b_2 a_2^* \dots a_m^* + \dots + b_m a_m^*)^*$$

This is the basis for an efficient algorithm for the exhaustive generation of m-compositions (Gray code), as described in:

E. Grazzini, E. Munarini, M. Poneti, S. Rinaldi, *On the generation of mcompositions and m-partitions*, 2006 Other results:

#### ◊ Asymptotic expansion:

$$c_n^{(m)} \sim_n \frac{1}{2m(\sqrt[m]{2}-1)} \left(\frac{\sqrt[m]{2}}{\sqrt[m]{2}-1}\right)^n$$

 $\diamond$  *m*-compositions without zero rows:

$$f_n^{(m)} = \sum_{k=0}^m \binom{m}{k} (-1)^{n-k} c_n^{(k)}$$

$$f^{(m)}(x) = \sum_{k=0}^{m} {m \choose k} (-1)^{n-k} \frac{(1-x)^k}{2(1-x)^k - 1}$$

 $\diamond$  *m*-compositions with palindromic rows:

$$\sum_{n \ge 0} p_n^{(m)} x^n = \frac{(1+x)^m}{2(1-x^2)^m - 1}$$

 $\diamond$  enumeration of various kind of  $m\mathchar`-colored$  bargraphs

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