# Computer Algebra and Power Series with Positive Coefficients 

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## Problem Statement

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Consider its Taylor expansion

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For some $r$, obviously no.
For some $r$, obviously yes.
For some $r$, it seems that yes, but nobody knows how to prove this.

## Simple Examples

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This mess does not have a closed form.
A positivity proof for $a_{n, m}$ is not obvious. (Try it.)

## Hard Examples

Askey/Gasper (1977) proved that the power series

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has only positive coefficients.
These proofs are really complicated!

## Open problems

Gillis/Reznick/Zeilberger conjectured in 1982 that

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has nonnegative coefficients for all $r \geq 4$.

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These conjectures are still open.

Can The Computer Help?

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Such an algorithm could be used for solving Diophantine equations, because

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\operatorname{rat}\left(x_{1}, \ldots, x_{r}\right)=\sum_{n_{1}, \ldots, n_{r}} \operatorname{poly}\left(n_{1}, \ldots, n_{r}\right)^{2} x_{1}^{n_{1}} \cdots x_{r}^{n_{r}}
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Hence: No such algorithm exists.

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Open problem: Does there exist an algorithm that can decide for a given $\left(a_{n}\right)$ whether $a_{N}=0$ for some $N$ ?
Such an algorithm exists if and only if there exists an algorithm for deciding whether a given univariate rational function has only positive coefficients.

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- For determining useful recurrence equations satisfied by the coefficients
- For deriving and deciding (proving or disproving) sufficient conditions that, if true, imply deciding certain sufficient conditions

Computer Algebra Tools

## Guessing and Proving Recurrence Equations

Task: Find a linear (possibly multivariate) recurrence equation with polynomial coefficients for the series coefficients of a given rational function.

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(Example.)

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## Back to the Examples

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The coefficients $a_{n, m}$ satisfy

$$
\begin{aligned}
& \frac{3}{16}(m+n+2)(m+n+3) a_{n, m} \\
& -\frac{1}{8}(m+n+3)(m+7 n+13) a_{n+1, m} \\
& +\frac{1}{2}(n+2)(2 n+5) a_{n+2, m}=0
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For positivity of $a_{n, m}$ it would be sufficient to show

$$
\begin{aligned}
& \forall N, M, A_{0}, A_{1}, A_{2} \in \mathbb{R}: \\
& N \geq M \geq 0 \wedge A_{1}>0 \wedge A_{0}>0 \\
& \wedge \frac{3}{16}(M+N+2)(M+N+3) A_{0} \\
& -\frac{1}{8}(M+N+3)(M+7 N+13) A_{1} \\
& +\frac{1}{2}(N+2)(2 N+5) A_{2}=0 \\
& \Rightarrow A_{2}>0
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Collins's algorithm tells us that this is wrong.

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The coefficients $a_{n, m}$ also satisfy

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\begin{aligned}
& -\frac{9}{32}(m+n+2)(m+n+3) \\
& \times(m+n+4)(m+n+5)(2 m+2 n+9) a_{n, m} \\
& +\frac{1}{4}(m+n+4)(m+n+5)(2 m+2 n+7) \\
& \times\left(2 m^{2}+16 n m+35 m+2 n^{2}+35 n+57\right) a_{n+1, m+1} \\
& -2(m+2)(2 m+5)(n+2) \\
& \times(2 n+5)(2 m+2 n+5) a_{n+2, m+2}=0
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Collin's algorithm tells us that this is also wrong.

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But it can be shown that the $a_{n, m}$ are diagonally increasing:

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is true. This proves the positivity result.

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- Easy: Compute some large (but finite) number of coefficients.
- Better: Prove positivity for some infinite set of coefficients.
- Best case: Prove positivity for all coefficients.



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- $a_{n, m, k}$ satisfies nice (but lengthy) recurrence equations.


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- But none of them implies positivity.
- The best we obtained were positivity implying recurrence equations in $n$ and $m$ when $k=1,2, \ldots, 16$ is fixed.


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- $a_{n, m, k}$ satisfies nice (but lengthy) recurrence equations.
- But none of them implies positivity.
- The best we obtained were positivity implying recurrence equations in $n$ and $m$ when $k=1,2, \ldots, 16$ is fixed.
- Because for each fixed $k$, a shorter recurrence equations are available.


## Szegö's Example

$$
\frac{1}{1-x-y-z+\frac{2}{3}(x y+x z+y z)}=\sum_{n, m, k} a_{n, m, k} x^{n} y^{m} z^{k}
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- Here it is possible to apply the monotonicity-by-induction reasoning.


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C A D \\
\Longleftrightarrow
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\frac{1}{1-x-y-z-w+\frac{3}{4}(x y+x z+x w+y z+y w+z w)}
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Such results can probably be obtained for any fixed $k, l, u, v, w$.

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- Automated guessing delivers 10 linearly independent multivariate recurrence equations with linear coefficients.
- A positivity-asserting linear-combination can be found by making an ansatz and solving a quantifier elimination problem for the undetermined coefficients.


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- This way, it can also be shown that Szegö's result cannot be shown by a first-order linear positivity-asserting recurrence with linear coefficients.


## Askey-Gasper's Example

The conjectured generalization

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The monotonicity-by-induction reasoning is applicable to this sum.

New Challenges \& Conclusion

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Computer experiments and partial proofs suggest that the following rational functions have positive coefficients:

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We can give partial proofs for both, but no full proofs.
(The first one is easily proven, as pointed out by Armin Straub a few days ago. The second remains open so far.)

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## Conclusion

- It is difficult to decide whether all the Taylor coefficients of a rational function are positive.
- Standard tools from Computer Algebra (Recurrence Guessing and Cylindrical Decomposition) can contribute to this topic.
- For nontrivial examples, we could obtain partial proofs in this way.
- This extends the computational evidence in support of these conjectures far beyond what was available so far.

