A bijection between 2-triangulations and pairs of non-crossing Dyck paths

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Let \mathfrak{C}_n be a regular *n*-gon.

A triangulation of \mathfrak{C}_n is a subdivision of \mathfrak{C}_n into triangles, using diagonals that do not cross.



- Every triangulation of \mathfrak{C}_n has exactly n-3 diagonals.
- The number of triangulations of \mathfrak{C}_n is

$$C_{n-2} = \frac{1}{n-1} \binom{2n-4}{n-2},$$

where C_m is the *m*-th Catalan number.

Dyck paths

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If (m, m) is the final point, we call m the size of the path.

Let \mathcal{D}_m be the set of Dyck paths of size m. Then, $|\mathcal{D}_m| = C_m$.

A bijection between triangulations and Dyck paths



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- \checkmark draw an N step,
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If the form (i, j) with i < j. Draw an *E* step at the end. **Definition.** A *j*-crossing is a set of j diagonals where any two of them cross.



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We could have defined a triangulation of \mathfrak{C}_n as a maximal set of diagonals with no 2-crossings.

Definition. A k-triangulation of \mathfrak{C}_n is a maximal set of diagonals with no (k+1)-crossings.



a 2-triangulation

Facts about generalized triangulations

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Every k-triangulation of \mathfrak{C}_n has exactly k(n-2k-1) diagonals.

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Theorem (Jonsson). The number of k-triangulations of \mathfrak{C}_n is

$$\det(C_{n-i-j})_{i,j=1}^{k} = \begin{vmatrix} C_{n-2} & C_{n-3} & \dots & C_{n-k} & C_{n-k-1} \\ C_{n-3} & C_{n-4} & \dots & C_{n-k-1} & C_{n-k-2} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ C_{n-k-1} & C_{n-k-2} & \dots & C_{n-2k+1} & C_{n-2k} \end{vmatrix}$$

Theorem (Lindström, Gessel-Viennot). The number of k-tuples (P_1, P_2, \ldots, P_k) of Dyck paths of size n - 2k such that each P_i never goes below P_{i+1} is given by the same determinant $\det(C_{n-i-j})_{i,j=1}^k$.



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- $\ \, \bullet \ \, k=1 \ \longrightarrow \ \, \mathsf{known}$
- \checkmark $k = 2 \longrightarrow$ we will see it next
- $k \ge 3 \longrightarrow$ open problem

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We omit them for simplicity.

Now there are exactly 2n - 10 diagonals.

We will construct a bijection between 2-triangulations and pairs of non-crossing Dyck paths.

Given a 2-triangulation, first we give an algorithm to color half of the crosses blue and the other half red.



At each iteration, one cross will be colored red and another blue, and two blocks will be merged.



- \checkmark Let r be the largest index so that row r has a cross in block r.
- \checkmark Color blue the leftmost uncolored cross in block r.
- Merge blocks r-2 and r-1.
 (If r=2, we consider that block 1 disappears when it is merged with "block 0".)
- Color red the rightmost uncolored cross in the merged block.



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The bijection (part II): from colored crosses to paths



 $\begin{array}{l} \alpha_i := \ \# \ \text{blue} \ \text{crosses} \ \text{in column} \ i \\ \beta_i := \ \# \ \text{red} \ \text{crosses} \ \text{in column} \ i \end{array}$

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Define

$$P = NE^{\alpha_5}NE^{\alpha_6}\cdots NE^{\alpha_{n-1}}NE^{\alpha_n}E$$
$$Q = NE^{\beta_4}NE^{\beta_5}\cdots NE^{\beta_{n-2}}NE^{\beta_{n-1}}E$$

How the bijection is obtained:

- Construct a generating tree for 2-triangulations.
- Construct a generating tree for pairs of non-crossing Dyck paths.
- Give an isomorphism between the generating trees.



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Generating rule:

$$(d_1, d_2, \dots, d_s) \longrightarrow \begin{cases} (i, d_j - i + 1, d_{j+1} + 1, d_{j+2}, \dots, d_s) : 1 \le j \le s - 1, \ 0 \le i \le d_j \\ \cup \{ (i, d_s - i + 1) : 0 \le i \le d_s + 1 \}. \end{cases}$$

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A generating tree for pairs of non-crossing Dyck paths



The nodes at level ℓ represent pairs of paths of size $\ell + 1$.

The two generating trees are isomorphic

We can give generating rules for 2-triangulations and for pairs of Dyck paths that yield isomorphic generating trees.



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For k-triangulations, we also have *diagonal flips*:



Fact: If we remove any diagonal, there a unique way to put it back to get another k-triangulation.

Open problem: Is there a polytope whose vertices correspond to *k*-triangulations and whose edges correspond to diagonal flips?

It should be a simple polytope of dimension k(n - 2k - 1).

For example, for k = 2 and n = 7, the graph of diagonal flips is



It can be realized as a a cyclic polytope in dimension 4, whose 3-dimensional facets are



Is there an analogous bijection between *k*-triangulations and *k*-tuples of non-crossing Dyck paths, for $k \ge 3$?

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Partial progress:

The same idea of splitting columns can be used to construct a generating tree for k-triangulations.



However, it is not clear what is the corresponding operation to generate children of a k-tuple of Dyck paths that would give an isomorphic generating tree.