

From algebraic sets to monomial linear bases by means of combinatorial algorithms.

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1.1 Let \mathbf{N} be the monoid of non-negative integers. Denote by $\mathbf{i} := (i_1, \dots, i_n)$ an arbitrary element in the power \mathbf{N}^n . The usual order on \mathbf{N} , as well as the partial order it induces on \mathbf{N}^n , will be denoted by \leq .

Define an *n-dimensional Ferrers diagram* to be any ideal of the poset \mathbf{N}^n , i.e. any non-empty subset $\mathcal{F} \subseteq \mathbf{N}^n$ such that $\mathbf{j} < \mathbf{i} \in \mathcal{F} \implies \mathbf{j} \in \mathcal{F}$. An element $\mathbf{i} = (i_1, \dots, i_n) \notin \mathcal{F}$ is said to be a *co-minimal element* for the Ferrers diagram \mathcal{F} if it is a minimal element of the complementary filter $\mathbf{N}^n \setminus \mathcal{F}$, i.e. if $(i_1, \dots, i_{r-1}, i_r - 1, i_{r+1}, \dots, i_n) \in \mathcal{F}$ for each r such that $i_r \geq 1$. Of course, $\mathbf{i} \notin \mathcal{F}$ is a co-minimal element iff $\mathcal{F}' := \{\mathbf{i}\} \cup \mathcal{F}$ is a Ferrers diagram.

We will write \preceq for any *term-ordering* on \mathbf{N}^n , i.e. a linear ordering which is compatible with the monoid structure on \mathbf{N}^n :

$$0 \prec \mathbf{i} \quad \text{for every } \mathbf{i} \neq 0 \text{ in } \mathbf{N}^n$$

$$\mathbf{i} \preceq \mathbf{j} \implies \mathbf{i} + \mathbf{r} \preceq \mathbf{j} + \mathbf{r} \quad \text{for every } \mathbf{i}, \mathbf{j}, \mathbf{r} \in \mathbf{N}^n.$$

It is well known that any term-ordering on \mathbf{N}^n is also a well-ordering.

1.2 Let K be a field and let $X := \{x_1, x_2, \dots, x_n\}$ be a given set of indeterminates. Let us consider the usual polynomial algebra $K[X] := K[x_1, \dots, x_n]$. Denote by $M_X \subseteq K[X]$ the free abelian monoid on X . The elements of M_X (i.e. the monic monomials) will be called *terms* of $K[X]$ and denoted by $\mathbf{x}^{\mathbf{i}} := x_1^{i_1} \cdots x_n^{i_n}$ with $\mathbf{i} := (i_1, \dots, i_n) \in \mathbf{N}^n$. The orders \leq and \preceq on \mathbf{N}^n , as well as the notion of Ferrers diagram, extend to M_X in an obvious way.

An ideal J of the algebra $K[X]$ is said to be *cofinite* if

$$\text{codim}(J) := \dim(K[X]/J) < \infty$$

Given a finite set $\mathcal{P} := \{P_1, \dots, P_N\} \subseteq K^n$, the ideal

$$\mathfrak{S}(\mathcal{P}) := \{p \in K[X] \mid (\forall P \in \mathcal{P})(p(P) = 0)\}$$

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MB2. [Put the first i points of the list \mathcal{P} in \mathcal{Q} .]

Set $\mathcal{Q} \leftarrow \{P_j \mid 1 \leq j \leq i\}$.

MB3. [Find which coordinate of \mathbf{d}_{i+1} has to be changed, say $d_{i+1,s}$.]

Set $s \leftarrow \max\{k \geq 1 \mid \pi_{k-1}(P_j) = \pi_{k-1}(P_{i+1}), \text{ for some } P_j \in \mathcal{Q}\}$.

($s - 1$ is the length of the longest initial segment shared by P_{i+1} and some point $P_j \in \mathcal{Q}$. If $s > 1$, then in successive steps this decreases.)

MB4. [Find the points that determine the s -th coordinate of \mathbf{d}_{i+1} .]

Set $\mathcal{E} \leftarrow \{j \mid P_j \in \mathcal{Q}, \pi_{s-1}(P_j) = \pi_{s-1}(P_{i+1}), \pi^{s+1}(\mathbf{d}_j) = \pi^{s+1}(\mathbf{d}_{i+1})\}$.

(Indices of the points of \mathcal{Q} which have the first $s - 1$ coordinates equal to those of P_{i+1} and whose corresponding elements in \mathcal{F} have the $n - s$ rightmost coordinates equal to those of \mathbf{d}_{i+1} . \mathcal{E} is always not-empty.)

MB5. [Assign the value to the s -th coordinate of \mathbf{d}_{i+1} .]

Set $d_{i+1,s} \leftarrow (1 + \max\{d_{j,s} \mid j \in \mathcal{E}\})$.

MB6. [Did you determine the first coordinate of \mathbf{d}_{i+1} ?]

If $s > 1$

MB6.1. [Find the points that determine another coordinate of \mathbf{d}_{i+1} .]

Set $\mathcal{Q} \leftarrow \{P_j \mid 1 \leq j \leq i, \pi^s(\mathbf{d}_j) = \pi^s(\mathbf{d}_{i+1}) = (d_{i+1,s}, \dots, d_{i+1,n})\}$.

(Points of \mathcal{P} whose corresponding elements in \mathcal{F} have the $n - s + 1$ rightmost coordinates equal to those of \mathbf{d}_{i+1} .)

MB6.2. [Is \mathcal{Q} empty?]

If $\mathcal{Q} \neq \emptyset$, return to step **MB3**.

MB7. [Increase i]

Set $i \leftarrow i + 1$. If $i < N$, return to step **MB2**.

MB8. [Done] Terminate the algorithm.

We put: $\mathcal{MB}(\mathcal{P}) := \{\mathbf{d}_1, \dots, \mathbf{d}_N\}$; $\delta_{\mathcal{P}}: \mathcal{P} \rightarrow \mathcal{MB}(\mathcal{P})$, $P_i \mapsto \mathbf{d}_i$. One could think that $\mathcal{MB}(\mathcal{P})$ is ill-defined, that is it depends on the order which has been used for arranging the points P_1, \dots, P_N when starting **Algorithm MB** (i.e. on the list \mathcal{P}) rather than on the set $\mathcal{P} = \{P_1, \dots, P_N\}$ itself. Well, this is not true. In fact, it is possible to prove the following propositions.

Prop. 1 Let $\mathcal{F}' := (\mathbf{d}'_1, \dots, \mathbf{d}'_N)$ be the list associated with the list of points $\mathcal{P}' = (P_{\sigma(1)}, \dots, P_{\sigma(N)})$, where $\sigma \in S_N$, by **Algorithm MB**. Then, for some $\tau \in S_N$, $(\mathbf{d}'_1, \dots, \mathbf{d}'_N) = (\mathbf{d}_{\tau(1)}, \dots, \mathbf{d}_{\tau(N)})$. \square

Corollary 1 If $\mathcal{P} \subseteq \mathcal{Q}$, then $\mathcal{MB}(\mathcal{P}) \subseteq \mathcal{MB}(\mathcal{Q})$. \square

of the algebraic set \mathcal{P} is cofinite. More generally, when K is algebraically closed, from the *Nullstellensatz* it follows that the ideal $J \subseteq K[X]$ is a cofinite ideal iff the algebraic set of J , i.e.

$$\mathcal{V}(J) := \{P \in K^n \mid (\forall p \in J)(p(P) = 0)\},$$

is finite. In this case we have $\#\mathcal{V}(J) \leq \text{codim}(J)$ (cfr. [3] p.23).

For a given ideal J , any linear basis \mathcal{L}_J of the quotient algebra $K[X]/J$ whose elements are of the form $[\mathbf{x}^{\mathbf{i}}]_J := \mathbf{x}^{\mathbf{i}} + J$ will be called a *monomial basis*. If $\mathcal{L}_J = \{\mathbf{x}^{\mathbf{i}} + J \mid \mathbf{i} \in L \subseteq \mathbb{N}^n\}$ is a monomial basis, we shall say that $\{\mathbf{x}^{\mathbf{i}} \mid \mathbf{i} \in L\}$ is a *system of representatives* for the monomial basis \mathcal{L}_J . Obviously, if $\mathcal{L}_J = \{\mathbf{x}^{\mathbf{i}} + J \mid \mathbf{i} \in L\}$ is a monomial basis of $K[X]/J$, then any polynomial $p \in K[X]$ is congruent modulo J to exactly one polynomial of the form $\sum_{\mathbf{i} \in L} a_{\mathbf{i}} \mathbf{x}^{\mathbf{i}}$, $a_{\mathbf{i}} \in K$. Given an arbitrary term-ordering \preceq on M_X , a monomial basis $\mathcal{L}_J = \{\mathbf{x}^{\mathbf{i}_1} + J, \dots, \mathbf{x}^{\mathbf{i}_N} + J\}$ with $\mathbf{x}^{\mathbf{i}_1} \prec \dots \prec \mathbf{x}^{\mathbf{i}_N}$ is said to be *minimal* with respect to \preceq if for any other monomial basis $\mathcal{L}'_J = \{\mathbf{x}^{\mathbf{i}'_1} + J, \dots, \mathbf{x}^{\mathbf{i}'_N} + J\}$ with $\mathbf{x}^{\mathbf{i}'_1} \prec \dots \prec \mathbf{x}^{\mathbf{i}'_N}$ we have $\mathbf{x}^{\mathbf{i}_j} \preceq \mathbf{x}^{\mathbf{i}'_j}$ for $j = 1, \dots, N$. Of course, both \mathcal{L}_J and \mathcal{L}'_J , $\mathcal{L}_J \neq \mathcal{L}'_J$, could be minimal with respect to different term-orderings. It is not hard to prove that if $\mathcal{L}_J = \{\mathbf{x}^{\mathbf{i}} + J \mid \mathbf{i} \in L\}$ is a minimal monomial basis then $L \subseteq \mathbb{N}^n$ is an n -dimensional Ferrers diagram.

1.3 In the search for a minimal monomial basis $\mathcal{L}_{\mathcal{P}}$, we present a purely combinatorial algorithm to get it from \mathcal{P} . In fact, making use of the **Algorithm MB** below, we associate a Ferrers diagram $\mathcal{MB}(\mathcal{P}) = \{\mathbf{d}_1, \dots, \mathbf{d}_N\} \subseteq \mathbb{N}^n$ to any finite set $\mathcal{P} := \{P_1, \dots, P_N\} \subseteq K^n$. The Ferrers diagram $\mathcal{MB}(\mathcal{P})$ gives the monomial basis $\mathcal{L}_{\mathcal{P}} = \{\mathbf{x}^{\mathbf{d}} + \mathfrak{S}(\mathcal{P}) \mid \mathbf{d} \in \mathcal{MB}(\mathcal{P})\}$ which is minimal with respect to the inverse lexicographical ordering $\preceq_{i.l.}$ with $x_1 \prec_{i.l.} x_2 \prec_{i.l.} \dots \prec_{i.l.} x_n$.

Put:

$$\begin{aligned} \mathcal{P} &:= \{P_1, \dots, P_N\} \subseteq K^n \\ \underline{\mathcal{P}} &:= (P_1, \dots, P_N) \\ \mathbf{d}_j &= (d_{j,1}, \dots, d_{j,n}) \in \mathbb{N}^n \\ \pi_s: K^n &\longrightarrow K^s, \quad (x_1, \dots, x_n) \mapsto (x_1, \dots, x_s) \\ (\pi_0(P) &\text{ is assumed to be the empty sequence.}) \\ \pi^s: K^n &\longrightarrow K^{n-s+1}, \quad (x_1, \dots, x_n) \mapsto (x_s, \dots, x_n) \end{aligned}$$

ALGORITHM MB. Given a list $\underline{\mathcal{P}} := (P_1, \dots, P_N)$ of points, we determine an ordered Ferrers diagram $\underline{\mathcal{F}} := (\mathbf{d}_1, \dots, \mathbf{d}_N)$.

MB1. [Initialize.]

Set $\mathbf{d}_1 \leftarrow \mathbf{d}_2 \leftarrow \dots \mathbf{d}_N \leftarrow (0, \dots, 0)$; $i \leftarrow 1$.

(In the beginning the coordinates of all the elements of $\underline{\mathcal{F}}$ are zero.)

Prop. 2 *It is possible to arrange the points P_1, \dots, P_N in a suitable list $\underline{\mathcal{P}}' = (P_{\sigma(1)}, \dots, P_{\sigma(N)})$ such that the elements in the corresponding list $\underline{\mathcal{F}}' = (\mathbf{d}_{\tau(1)}, \dots, \mathbf{d}_{\tau(N)})$ are arranged according to the inverse lexicographical order $\preceq_{i.l.}$: $s < t \Rightarrow \mathbf{d}_{\tau(s)} \prec_{i.l.} \mathbf{d}_{\tau(t)}$ \square*

Lemma 1 *Let $\underline{\mathcal{F}} := (\mathbf{d}_1, \dots, \mathbf{d}_N)$ be the Ferrers diagram associated to $\underline{\mathcal{P}} := (P_1, \dots, P_N)$ by Algorithm MB; let $\mathbf{d}_N = (d_{N,1}, \dots, d_{N,n})$. If $d_{N,i} \neq 0$, then there is some $k < N$ such that $\mathbf{d}_k = (d_{N,1}, \dots, d_{N,i-1}, d_{N,i} - 1, d_{N,i+1}, \dots, d_{N,n})$.*

Proof. In order to calculate the first $i - 1$ coordinates $d_{N,1}, \dots, d_{N,i-1}$ of $\mathbf{d}_N = \delta_{\underline{\mathcal{P}}}(P_N)$, we have to consider the set $\{P_{j_1}, \dots, P_{j_s}, P_{j_{s+1}} = P_N\} \subseteq \mathcal{P}$ $j_1 < \dots < j_s < N$, of all points $P_{j_r} \in \mathcal{P}$ such that $\pi^i(\delta_{\underline{\mathcal{P}}}(P_{j_r})) = (d_{N,i}, \dots, d_{N,n})$. Putting $\tilde{\mathcal{P}} := \{\pi_{i-1}(P_{j_1}), \dots, \pi_{i-1}(P_{j_s}), \pi_{i-1}(P_N)\}$, we have $\pi_{i-1}(\delta_{\underline{\mathcal{P}}}(P_N)) = \delta_{\tilde{\mathcal{P}}}(\pi_{i-1}(P_N)) = (d_{N,1}, \dots, d_{N,i-1})$. For every $r \in \{1, \dots, s+1\}$, there is a point $P_{j'_r} \in \mathcal{P}$, $j'_r < j_r$, such that $d_{j'_r,i} = d_{j_r,i} - 1 = d_{N,i} - 1$, $d_{j'_r,i+1} = d_{j_r,i+1} = d_{N,i+1}, \dots, d_{j'_r,n} = d_{j_r,n} = d_{N,n}$ and $\pi_{i-1}(P_{j'_r}) = \pi_{i-1}(P_{j_r})$. Up to a suitable rearrangement of the points in $\underline{\mathcal{P}}$, we may assume that $j'_1 < \dots < j'_{s+1}$. When calculating $\pi_{i-1}(\delta_{\underline{\mathcal{P}}}(P_{j'_{s+1}}))$ we have to consider the set $\mathcal{Q} := \{P_l \in \mathcal{P} \mid l \leq j'_{s+1}, \pi^i(\mathbf{d}_l) = (d_{j'_{s+1},i}, \dots, d_{j'_{s+1},n}) = (d_{N,i} - 1, d_{N,i+1}, \dots, d_{N,n})\}$. Of course $P_{j'_r} \in \mathcal{Q}$. Let $\tilde{\mathcal{Q}} := \{\pi_{i-1}(P) \mid P \in \mathcal{Q}\} \subseteq K^{i-1}$; we have $\tilde{\mathcal{P}} \subseteq \tilde{\mathcal{Q}}$. Hence $\mathcal{MB}(\tilde{\mathcal{P}}) \subseteq \mathcal{MB}(\tilde{\mathcal{Q}}) \subseteq N^{i-1}$; in particular $(d_{N,1}, \dots, d_{N,i-1}) \in \mathcal{MB}(\tilde{\mathcal{Q}})$. It follows that there exists a point $P_k \in \mathcal{Q}$ such that $\delta_{\underline{\mathcal{P}}}(P_k) = (d_{N,1}, \dots, d_{N,i-1}, d_{N,i} - 1, d_{N,i+1}, \dots, d_{N,n})$. \square

As a straightforward consequence of Lemma 1 we get:

Prop. 3 *The set $\mathcal{MB}(\mathcal{P})$ is an n -dimensional Ferrers diagram.* \square

Prop. 4 *The set $\mathcal{L}_{\mathcal{P}} = \{\mathbf{x}^{\mathbf{d}} + \mathfrak{S}(\mathcal{P}) \mid \mathbf{d} \in \mathcal{MB}(\mathcal{P})\}$ is a monomial linear basis for $K[X]/\mathfrak{S}(\mathcal{P})$.*

Proof. By induction on the dimension n of K^n .

For $n = 1$, we have $\mathcal{P} = \{\rho_1, \dots, \rho_n\}$, $\rho_i \in K$, and $\mathfrak{S}(\mathcal{P}) = (g)$, with $g = \prod_{i=1}^N (x - \rho_i)$. Algorithm MB gives $\mathcal{MB}(\mathcal{P}) = \{0, 1, \dots, N - 1\}$; hence $\mathcal{L}_{\mathcal{P}} = \{\mathbf{x}^{\mathbf{d}} + \mathfrak{S}(\mathcal{P}) \mid \mathbf{d} \in \mathcal{MB}(\mathcal{P})\} = \{1 + (g), x + (g), \dots, x^{N-1} + (g)\}$, which is a minimal monomial basis for $K[X]/(g)$.

Suppose now that the statement is true for every finite subset of $K^{n'}$, $n' < n$, and prove it for $\mathcal{P} = \{P_1, \dots, P_N\} \subset K^n$. As $\dim K[X]/\mathfrak{S}(\mathcal{P}) = N = \#\mathcal{MB}(\mathcal{P})$, it remains to prove that the residue classes mod. $\mathfrak{S}(\mathcal{P})$ of the monomials $\mathbf{x}^{\mathbf{d}}$, $\mathbf{d} \in \mathcal{MB}(\mathcal{P})$, are linearly independent over $K[X]/\mathfrak{S}(\mathcal{P})$; in other words, we have to prove that any polynomial of the form

$$(1) \quad p(x_1, \dots, x_n) = \sum_{\mathbf{d} \in \mathcal{MB}(\mathcal{P})} \alpha_{\mathbf{d}} \mathbf{x}^{\mathbf{d}} \in \mathfrak{S}(\mathcal{P}), \quad \alpha_{\mathbf{d}} \in K$$

is the zero polynomial.

Putting $D := \mathcal{MB}(\mathcal{P})$, $D_r := \{d = (d_1, \dots, d_n) \in D \mid d_n = r\}$ and $\mathcal{P}_r := \{P \in \mathcal{P} \mid \delta_{\mathcal{P}}(P) \in D_r\}$, it is easy to check that

$$(2) \quad \mathcal{MB}(\pi_{n-1}(\mathcal{P}_r)) = \pi_{n-1}(D_r).$$

Let us write down polynomial (1) in the form

$$(3) \quad p(x_1, \dots, x_n) = \sum_{r=0}^h p_r(x_1, \dots, x_{n-1})x_n^r$$

where $h = \max\{d_n \mid d = (d_1, \dots, d_n) \in D\}$ and

$$(4) \quad p_r(x_1, \dots, x_{n-1}) = \sum_{(d_1, \dots, d_{n-1}) \in \pi_{n-1}(D_r)} \alpha_{(d_1, \dots, d_{n-1}, r)} x_1^{d_1} \cdots x_{n-1}^{d_{n-1}}.$$

The polynomial $p(x_1, \dots, x_n) \in \mathfrak{S}(\mathcal{P})$ has to vanish at every point in \mathcal{P} . Consider a point $P = (a_1, \dots, a_n) \in \mathcal{P}_h \subseteq \mathcal{P}$; there exist in \mathcal{P} exactly $h+1$ points that have the same first $n-1$ coordinates as $P = (a_1, \dots, a_n)$. It follows that the polynomial

$$p(a_1, \dots, a_{n-1}, x_n) = \sum_{r=0}^h p_r(a_1, \dots, a_{n-1})x_n^r$$

vanishes identically. In particular $p_h(a_1, \dots, a_{n-1}) = 0$ for every $(a_1, \dots, a_{n-1}) \in \mathcal{Q} := \pi_{n-1}(\mathcal{P}_h)$. Hence

$$(5) \quad p_h(x_1, \dots, x_{n-1}) \in \mathfrak{S}(\mathcal{Q}) \subseteq K[x_1, \dots, x_{n-1}].$$

By (2) and (4) we have

$$(6) \quad p_h(x_1, \dots, x_{n-1}) = \sum_{(d_1, \dots, d_{n-1}) \in \mathcal{MB}(\mathcal{Q})} \alpha_{(d_1, \dots, d_{n-1}, h)} x_1^{d_1} \cdots x_{n-1}^{d_{n-1}}$$

Because of the inductive hypothesis, the set $\{x_1^{d_1} \cdots x_{n-1}^{d_{n-1}} + \mathfrak{S}(\mathcal{Q}) \mid (d_1, \dots, d_{n-1}) \in \mathcal{MB}(\mathcal{Q})\}$ is a monomial basis for $K[x_1, \dots, x_{n-1}]/\mathfrak{S}(\mathcal{Q})$. From this and from (5) we deduce that polynomial (6) vanishes identically. Hence

$$(7) \quad p(x_1, \dots, x_n) = \sum_{r=0}^{h-1} p_r(x_1, \dots, x_{n-1})x_n^r$$

Arguing for $r = h-1, h-2, \dots, 1, 0$ as for $r = h$, we conclude that $p_r(x_1, \dots, x_{n-1})$ is the zero polynomial for every $r \in \{0, \dots, h\}$. \square

Prop. 5 *The monomial linear basis $\mathcal{L}_{\mathcal{P}} = \{x^d + \mathfrak{S}(\mathcal{P}) \mid d \in \mathcal{MB}(\mathcal{P})\}$ for $K[X]/\mathfrak{S}(\mathcal{P})$ is minimal with respect to the inverse lexicographical order $\preceq_{i.1}$ on $M_{\mathcal{X}}$.*

Proof. Let $\mathcal{P} = (P_1, \dots, P_N)$, $\underline{d} := \mathcal{MB}(\mathcal{P}) = (\mathbf{d}_1, \dots, \mathbf{d}_N)$; $\mathbf{d}_i = (d_{i,1}, \dots, d_{i,n})$, $P_i = (a_{i,1}, \dots, a_{i,n})$ for $i = 1, \dots, N$. Let $h := \max\{d_{i,n} \mid i = 1, \dots, N\}$, and $\mathcal{P}_h = \{P_{j_1}, \dots, P_{j_m}\} \subseteq \mathcal{P}$; for any P_{j_r} there are in \mathcal{P} exactly $h + 1$ points which have the same $n - 1$ coordinates as P_{j_r} ; let us denote them by $Q_{j_r,0}, \dots, Q_{j_r,h} = P_{j_r}$. Up to a suitable rearrangement we may assume that they are the last $(h + 1)m$ elements in the list \mathcal{P} , i.e.

$$(Q_{j_1,0}, \dots, Q_{j_1,h}, \dots, Q_{j_m,0}, \dots, Q_{j_m,h}) = (P_{N-(h+1)m+1}, P_{N-(h+1)m+2}, \dots, P_N),$$

so that

$$d_{N-(h+1)(m-1),n} = d_{N-(h+1)(m-2),n} = \dots = d_{N-(h+1),n} = d_{N,n} = h$$

We have to prove that for any $\mathbf{d}' = (d'_1, \dots, d'_n)$ such that $\mathbf{d}' \prec_{i,l} \mathbf{d}_N$ there exists in $\mathfrak{S}(\mathcal{P})$ a polynomial of the form

$$(8) \quad \sum_{i=1}^{N-1} \alpha_i \mathbf{x}^{\mathbf{d}_i} + \alpha \mathbf{x}^{\mathbf{d}'} \in \mathfrak{S}(\mathcal{P}).$$

Because of MB.6.1 of Algorithm MB, without loss of generality we may assume that $d'_n < d_{N,n}$. Observe that (8) is equivalent to

$$(9) \quad \sum_{i=1}^{N-1} \alpha_i a_{s,1}^{d_{i,1}} \dots a_{s,n}^{d_{i,n}} + \alpha a_{s,1}^{d'_1} \dots a_{s,n}^{d'_n} = 0, \quad P_s = (a_{s,1}, \dots, a_{s,n}) \in \mathcal{P}$$

Hence, it is enough to prove that the N by N matrix A whose $s - th$ row is

$$a_{s,1}^{d_{i,1}} \dots a_{s,n}^{d_{i,n}} \quad \dots \quad a_{s,1}^{d_{N-1,1}} \dots a_{s,n}^{d_{N-1,n}} \quad a_{s,1}^{d'_1} \dots a_{s,n}^{d'_n}$$

(that is, the evaluation at P_s of the list of monomials $\mathbf{x}^{\mathbf{d}_1}, \dots, \mathbf{x}^{\mathbf{d}_{N-1}}, \mathbf{x}^{\mathbf{d}'}$) is a singular matrix. We shall prove this by showing that the submatrix A' consisting of the last $(h + 1)m$ rows of A has rank less than $(h + 1)m$.

Let X be any minor of order $(h + 1)m$ of A' ; notice that X can be given the form

$$(10) \quad X = \sum M_1 \cdot M_2 \cdot \dots \cdot M_m$$

where M_i is a minor of A' which consists of the $h + 1$ rows whose indices are $N - (h + 1)(m - i + 1) + 1, N - (h + 1)(m - i + 1) + 2, \dots, N - (h + 1)(m - i)$. Since all the points $P_{N-(h+1)(m-i+1)+1}, P_{N-(h+1)(m-i+1)+2}, \dots, P_{N-(h+1)(m-i)}$ have the same $(n - 1)$ first coordinates, the minor M_i is different from zero only if its $(h + 1)$ columns correspond to monomials $\mathbf{x}^{\mathbf{i}} = x_1^{i_1} \dots x_{n-1}^{i_{n-1}} x_n^{i_n}$ such that their exponents i_n 's have all the possible values $0, 1, \dots, h$. On the other hand, there are no more than $m - 1$ such $(h + 1)$ -tuples of different columns (since only $m - 1$ among the monomials $\mathbf{x}^{\mathbf{d}_1}, \dots, \mathbf{x}^{\mathbf{d}_{N-1}}, \mathbf{x}^{\mathbf{d}'}$ have h as last exponent);

it follows that at least one of the m minors M_i 's in (10) is zero. Hence $X = 0$.
□

1.4 In the case where $n = 2$ **Algorithm MB** can be given the following simplified form.

Assume that

$$\mathcal{P} = \{(a_1, b_{11}), \dots, (a_1, b_{1h_1}), \dots, (a_m, b_{m1}), \dots, (a_m, b_{mh_m})\}$$

with $h_1 + \dots + h_m = N$ and $i \neq j \Rightarrow a_i \neq a_j$. Then,

$$\mathcal{MB}(\mathcal{P}) = \{(p, q) \mid 0 \leq p < m, 0 \leq q < h_p\}.$$

1.5 All that we have stated up to now can be generalized to what we might call *algebraic multisets* in the following sense.

Consider the linear map

$$\begin{aligned} \mathbf{D}_{\mathbf{i}}: K[X] &\longrightarrow K[X] \\ \mathbf{x}^{\mathbf{h}} &\longmapsto \binom{\mathbf{h}}{\mathbf{i}} \mathbf{x}^{\mathbf{h}-\mathbf{i}} \end{aligned}$$

where $\mathbf{h} = (h_1, \dots, h_n)$, $\mathbf{i} = (i_1, \dots, i_n) \in \mathbf{N}$ and $\binom{\mathbf{h}}{\mathbf{i}} := \binom{h_1}{i_1} \cdots \binom{h_n}{i_n}$. Observe that when the field K has characteristic zero, then

$$\mathbf{D}_{\mathbf{i}} = \frac{1}{\mathbf{i}!} \mathbf{D}^{\mathbf{i}} := \frac{1}{\mathbf{i}!} \frac{\partial^{i_1 + \dots + i_n}}{\partial x_1^{i_1} \cdots \partial x_n^{i_n}}$$

where $\mathbf{i} = (i_1, \dots, i_n)$ and $\mathbf{i}! = i_1! \cdots i_n!$.

Let v_P be the evaluation map at the point P :

$$\begin{aligned} v_P: K[X] &\longrightarrow K \\ q &\longmapsto q(P) \end{aligned}$$

Define the linear map $v_P^{\mathbf{i}}$ as the composition $v_P \circ \mathbf{D}_{\mathbf{i}}$, i.e.

$$\begin{aligned} v_P^{\mathbf{i}}: K[X] &\longrightarrow K \\ q &\longmapsto (\mathbf{D}_{\mathbf{i}}q)(P). \end{aligned}$$

For every ideal J of $K[x_1, \dots, x_n]$ and every $P \in \mathcal{V}(J)$, put

$$\mathcal{F}_J(P) := \left\{ \mathbf{i} \in \mathbf{N}^n \mid (\forall p \in J) \left(v_P^{\mathbf{i}}(p) = 0 \right) \right\}$$

Prop. 6 (i) $\mathcal{F}_J(P)$ is a Ferrers diagram; (ii) if (g_1, \dots, g_s) is a system of generators of J , then $\mathcal{F}_J(P)$ is the largest Ferrers diagram contained in the set

$$\left\{ \mathbf{i} \in \mathbb{N}^n \mid v_{\mathbf{P}}^{\mathbf{i}}(g_j) = 0 \text{ for every } 1 \leq j \leq s \right\}.$$

We define a *finite n -dimensional algebraic multiset*, or simply an *algebraic multiset*, to be a set $\wp := \{(P_1, \mathcal{F}_1), \dots, (P_N, \mathcal{F}_N)\}$; each element $(P_j, \mathcal{F}_j) \in \wp$ consists of a point P_j of K^n together with a Ferrers diagram $\mathcal{F}_j \subset \mathbb{N}^n$, which will be called the *algebraic diagram* of the point P_j . We shall freely make use of the notation $(P, \mathbf{i}) \in \wp$, or also $P \in \wp$, to mean that for some $j \in \{1, \dots, N\}$, $P = P_j$ and $\mathbf{i} \in \mathcal{F}_j$. With every algebraic multiset $\wp = \{(P_1, \mathcal{F}_1), \dots, (P_N, \mathcal{F}_N)\}$ we associate the set

$$\mathfrak{S}(\wp) := \left\{ p \in K[X] \mid (\forall j)(\forall \mathbf{i} \in \mathcal{F}_j) \left(v_{P_j}^{\mathbf{i}}(p) = 0 \right) \right\}.$$

It is not difficult to prove that $\mathfrak{S}(\wp)$ is a cofinite ideal of $K[X]$ and that $\mathcal{F}_{\mathfrak{S}(\wp)}(P_j) = \mathcal{F}_j$ for every $j \in \{1, \dots, N\}$. Moreover, one can prove that

$$\text{codim } \mathfrak{S}(\wp) = \#\wp := \sum_{j=1}^N \#\mathcal{F}_j.$$

The question now is: how do we get a monomial linear basis for $K[X]/\mathfrak{S}(\wp)$? Once more the problem can be solved by applying a slightly modified version of **Algorithm MB** (in fact **Algorithm MB** itself with a few obvious changes) to a suitable set $\mathcal{R}(\wp) \subset (K \times \mathbb{N})^n$ associated with the algebraic multiset \wp . $\mathcal{R}(\wp)$ will be called the *umbral representation* of \wp . To be precise, consider the bijection

$$\begin{aligned} u: K^n \times \mathbb{N}^n &\longrightarrow (K \times \mathbb{N})^n \\ ((a_1, \dots, a_n), (i_1, \dots, i_n)) &\longmapsto ((a_1, i_1), \dots, (a_n, i_n)). \end{aligned}$$

Put

$$\mathcal{R} = \mathcal{R}(\wp) := \{u(P, \mathbf{i}) \mid (P, \mathbf{i}) \in \wp\}.$$

and

$$\mathcal{MB}(\wp) := \mathcal{MB}(\mathcal{R}).$$

(the symbol \mathcal{MB} on the right-hand side represents the operator defined by **Algorithm MB**). It is possible to prove that $\mathcal{MB}(\wp)$ satisfies properties analogous to those of $\mathcal{MB}(\mathcal{P})$; in particular, (i) $\mathcal{MB}(\wp)$ is a Ferrers diagram and (ii) the set $B := \{\mathbf{x}^{\mathbf{i}} + \mathfrak{S}(\wp) \mid \mathbf{i} \in \mathcal{MB}(\wp)\}$ is a monomial linear basis of $K[X]/\mathfrak{S}(\wp)$.

1.6 The above algorithms may come in handy for solving various problems. Let us examine a few of them.

I. First of all, let us see how to determine a system of generators $(\gamma_1, \dots, \gamma_r)$ for the ideal $\mathfrak{S}(\wp)$ of a finite algebraic multiset $\wp := \{(P_1, \mathcal{F}_1), \dots, (P_N, \mathcal{F}_N)\}$. In fact, the set $\{\gamma_1, \dots, \gamma_r\}$ we shall obtain is also a *reduced Gröbner basis* of $\mathfrak{S}(\wp)$. It goes without saying that the same procedure works also when \wp is a finite algebraic set.

Let B be the monomial linear basis obtained in 1.5 and let $Y = \{\mathbf{x}_{r_1}, \dots, \mathbf{x}_{r_r}\} \subset M_X$ be the minimal set of terms such that $B = M_X \setminus \sum_{i=1}^r \mathbf{x}_{r_i} \cdot M_X$. For each $\mathbf{x}_{r_i} \in Y$ determine a polynomial $\gamma_i \in K[X]$ in the form of a determinant in the following way. The first row of γ_i is the list $(\mathbf{x}_{d_1}, \dots, \mathbf{x}_{d_m}, \mathbf{x}_{r_i})$ where $\mathbf{x}_{d_j} \in B$ and $m = \#B = \text{codim}(\mathfrak{S}(\wp))$; the successive rows are lists of the form $(v_{P_j}^i(\mathbf{x}_{d_1}), \dots, v_{P_j}^i(\mathbf{x}_{d_m}), v_{P_j}^i(\mathbf{x}_{r_i}))$, one for each $(P_j, i) \in \wp$. It can be proved that the list $(\gamma_1, \dots, \gamma_r)$ is a reduced Gröbner basis of $\mathfrak{S}(\wp)$.

II. Consider a linear form $f \in K[X]^* \cong K[[X]]$. If $\text{Ker}(f)$ contains a cofinite ideal J of $K[X]$, then f is said to be an *n-linearly recursive function* and J is called a *characteristic ideal* of f . This notion has been introduced in [4] as a generalization of that of *linearly recursive sequence*, to which it reduces when $n = 1$. n-linearly recursive functions may also be regarded as elements of the dual bialgebra of the usual polynomial bialgebra on $K[X]$. When working on these subjects, it may happen that examples (perhaps, *suitable examples*) of n-linearly recursive functions are needed. How to construct them? It is convenient to divide the answer to this question into two parts.

(A) Let us first suppose that we know a system of generators (g_1, \dots, g_s) of the characteristic ideal J of the n-linearly recursive functions we are considering. In this case we may calculate a reduced Gröbner basis $G_{\preceq} := \text{RGB}(g_1, \dots, g_s)$ of (g_1, \dots, g_s) with respect to some term-ordering \preceq on M_X . Let $G_{\preceq} = (\gamma_1, \dots, \gamma_r)$, $\gamma_j \in K[X]$, and let $\xi_j \in M_X$ be the leading term (with respect to \preceq) of the polynomial γ_j . Then the set $B = M_X \setminus \sum_{j=1}^r \xi_j \cdot M_X$ is a monomial linear basis for $K[X]/J$, which is minimal with respect to \preceq . It follows that any n-linearly recursive function whose characteristic polynomial is J is uniquely determined by the set of its *initial values* $\{f(\mathbf{x}^d) \mid \mathbf{x}^d \in B\}$, all the other values $f(\mathbf{x}^i)$, $\mathbf{x}^i \notin B$, being calculated by making use of the polynomials $\gamma_j \in G_{\preceq}$ as *scales of recurrence*.

(B) If instead the characteristic ideal is not given, we are quite unlikely to obtain one of them (remember it must be cofinite!) just choosing at random a set of generators: most of the times we would get a non-cofinite ideal or, when cofinite, a trivial one. The correct way to do this consists instead in choosing a finite algebraic multiset (possibly, a finite algebraic set) \wp and then determining both the monomial linear basis B of $K[X]/\mathfrak{S}(\wp)$ and a set of generators for $\mathfrak{S}(\wp)$ by means of the machinery described in the previous sections.

III. Lastly, consider the following interpolation problem: given a finite n-dimensional algebraic multiset $\wp := \{(P_1, \mathcal{F}_1), \dots, (P_N, \mathcal{F}_N)\}$ and a set of values $\{\alpha_{j,i} \mid j = 1, \dots, N \text{ and } i \in \mathcal{F}_j\} \subset K$ determine the *unique* polynomial p of the

form $\sum_{\mathbf{x}_i \in B} a_i \mathbf{x}_i^{\mathbf{i}}$ ($a_i \in K$ and B is a monomial linear basis of $K[X]/\mathfrak{S}(\wp)$) for which we have $v_{P_j}^{\mathbf{i}}(p) = \alpha_{j,\mathbf{i}}$.

This problem appears to be an n -dimensional natural generalization of the unidimensional one which is solved by means of Lagrange interpolation formula (though a thorough analysis of these two shows that in some respects the analogy necessarily fails). Once more, the key point for solving this problem is to determine the monomial basis B . Let $B = \{\mathbf{x}_{\mathbf{d}_1}, \dots, \mathbf{x}_{\mathbf{d}_m}\}$ and let \mathcal{A} be the $m \times m$ matrix whose rows are of the form $(v_{P_j}^{\mathbf{i}}(\mathbf{x}_{\mathbf{d}_1}), \dots, v_{P_j}^{\mathbf{i}}(\mathbf{x}_{\mathbf{d}_m}))$, one for each $(P_j, \mathbf{i}) \in \wp$. Consider the vector $\bar{\alpha}$ whose components are the values $\alpha_{j,\mathbf{i}} = v_{P_j}^{\mathbf{i}}(p)$ (arranged according to the order that has been used for the rows of \mathcal{A}). The components of the vector $\bar{\beta} := \mathcal{A}^{-1} \cdot \bar{\alpha}$ are the coefficients of the desired polynomial.

Appendix.

1) Example of Algorithm MB for a 4-dimensional set.

\mathcal{P}	\longleftrightarrow	\mathcal{F}
$P_1 = (2, 3, 9, 4)$	\longleftrightarrow	$d_1 = (0, 0, 0, 0)$
$P_2 = (2, 5, 7, 3)$	\longleftrightarrow	$d_2 = (0, 1, 0, 0)$
$P_3 = (2, 3, 3, 2)$	\longleftrightarrow	$d_3 = (0, 0, 1, 0)$
$P_4 = (2, 5, 5, 1)$	\longleftrightarrow	$d_4 = (0, 1, 1, 0)$
$P_5 = (6, 1, 1, 3)$	\longleftrightarrow	$d_5 = (1, 0, 0, 0)$
$P_6 = (2, 3, 3, 6)$	\longleftrightarrow	$d_6 = (0, 0, 0, 1)$
$P_7 = (8, 3, 4, 0)$	\longleftrightarrow	$d_7 = (2, 0, 0, 0)$
$P_8 = (6, 5, 6, 3)$	\longleftrightarrow	$d_8 = (1, 1, 0, 0)$
$P_9 = (4, 7, 6, 6)$	\longleftrightarrow	$d_9 = (3, 0, 0, 0)$
$P_{10} = (4, 1, 7, 7)$	\longleftrightarrow	$d_{10} = (2, 1, 0, 0)$
$P_{11} = (2, 5, 7, 8)$	\longleftrightarrow	$d_{11} = (0, 1, 0, 1)$
$P_{12} = (4, 3, 0, 3)$	\longleftrightarrow	$d_{12} = (0, 2, 0, 0)$
$P_{13} = (1, 1, 2, 5)$	\longleftrightarrow	$d_{13} = (4, 0, 0, 0)$
$P_{14} = (2, 1, 9, 0)$	\longleftrightarrow	$d_{14} = (1, 2, 0, 0)$
$P_{15} = (4, 1, 7, 6)$	\longleftrightarrow	$d_{15} = (1, 0, 0, 1)$
$P_{16} = (8, 1, 8, 0)$	\longleftrightarrow	$d_{16} = (3, 1, 0, 0)$
$P_{17} = (2, 1, 7, 6)$	\longleftrightarrow	$d_{17} = (0, 2, 1, 0)$
$P_{18} = (4, 3, 3, 3)$	\longleftrightarrow	$d_{18} = (1, 0, 1, 0)$
$P_{19} = (4, 1, 1, 0)$	\longleftrightarrow	$d_{19} = (1, 1, 1, 0)$
$P_{20} = (8, 1, 2, 4)$	\longleftrightarrow	$d_{20} = (2, 0, 1, 0)$
$P_{21} = (6, 3, 4, 8)$	\longleftrightarrow	$d_{21} = (2, 2, 0, 0)$
$P_{22} = (2, 9, 6, 7)$	\longleftrightarrow	$d_{22} = (0, 3, 0, 0)$
$P_{23} = (2, 3, 7, 1)$	\longleftrightarrow	$d_{23} = (0, 0, 2, 0)$
$P_{24} = (2, 1, 7, 5)$	\longleftrightarrow	$d_{24} = (0, 2, 0, 1)$
$P_{25} = (4, 1, 1, 2)$	\longleftrightarrow	$d_{25} = (0, 0, 1, 1)$
$P_{26} = (2, 7, 6, 5)$	\longleftrightarrow	$d_{26} = (0, 4, 0, 0)$
$P_{27} = (2, 5, 5, 4)$	\longleftrightarrow	$d_{27} = (1, 0, 1, 1)$
$P_{28} = (2, 3, 7, 7)$	\longleftrightarrow	$d_{28} = (0, 1, 1, 1)$
$P_{29} = (2, 3, 0, 2)$	\longleftrightarrow	$d_{29} = (0, 0, 3, 0)$
$P_{30} = (2, 1, 1, 1)$	\longleftrightarrow	$d_{30} = (0, 1, 2, 0)$
$P_{31} = (4, 3, 3, 7)$	\longleftrightarrow	$d_{31} = (1, 1, 0, 1)$
$P_{32} = (8, 5, 6, 4)$	\longleftrightarrow	$d_{32} = (3, 2, 0, 0)$
$P_{33} = (4, 5, 5, 2)$	\longleftrightarrow	$d_{33} = (1, 3, 0, 0)$
$P_{34} = (2, 1, 1, 8)$	\longleftrightarrow	$d_{34} = (0, 2, 1, 1)$
$P_{35} = (8, 1, 1, 1)$	\longleftrightarrow	$d_{35} = (1, 0, 2, 0)$
$P_{36} = (4, 5, 5, 5)$	\longleftrightarrow	$d_{36} = (1, 2, 0, 1)$
$P_{37} = (6, 1, 8, 8)$	\longleftrightarrow	$d_{37} = (3, 0, 1, 0)$
$P_{38} = (2, 7, 8, 4)$	\longleftrightarrow	$d_{38} = (0, 3, 1, 0)$

The system of representatives for the corresponding monomial basis is given by

1	x_1	x_1^2	x_1^3	x_1^4	x_3	x_1x_3	$x_1^2x_3$	$x_1^3x_3$
x_2	x_1x_2	$x_1^2x_2$	$x_1^3x_2$		x_2x_3	$x_1x_2x_3$	$x_1^3x_3$	
x_2^2	$x_1x_2^2$	$x_1^2x_2^2$	$x_1^3x_2^2$		$x_2^2x_3$			
x_2^3	$x_1x_2^3$							
x_2^4								
x_3^2	$x_1x_3^2$				x_3^3			
$x_2x_3^2$								
x_4	x_1x_4				x_3x_4	$x_1x_3x_4$		
x_2x_4	$x_1x_2x_4$				$x_2x_3x_4$			
$x_2^2x_4$	$x_1x_2^2x_4$				$x_2^2x_3x_4$			

2) Example of Algorithm MB for a 3-dimensional algebraic multiset.

Consider the algebraic multiset \wp given by

$$\begin{aligned}
 P_1 = (0, 0, 0) \quad \mathcal{F}_1 &= \begin{bmatrix} (0, 0, 0) & (1, 0, 0) & & (0, 0, 1) \\ (0, 1, 0) & & & \end{bmatrix} \\
 P_2 = (0, 0, 1) \quad \mathcal{F}_2 &= \begin{bmatrix} (0, 0, 0) & (1, 0, 0) & (2, 0, 0) & (0, 0, 1) & (1, 0, 1) \end{bmatrix} \\
 P_3 = (1, 1, 0) \quad \mathcal{F}_3 &= \begin{bmatrix} (0, 0, 0) & (1, 0, 0) & & (0, 0, 1) & (1, 0, 1) & (0, 0, 2) \\ (0, 1, 0) & (1, 1, 0) & & (0, 1, 1) & & (0, 1, 2) \end{bmatrix} \\
 P_4 = (1, 1, 1) \quad \mathcal{F}_4 &= \begin{bmatrix} (0, 0, 0) \\ (0, 1, 0) \end{bmatrix}
 \end{aligned}$$

The umbral representation $\mathcal{R} = \mathcal{R}(\wp)$ of \wp (whose elements are intentionally arranged at random) as well as the corresponding diagram $\mathcal{MB}(\mathcal{R})$ are

\wp	\longleftrightarrow	\mathcal{R}	\longleftrightarrow	$\mathcal{MB}(\mathcal{R})$
$(1, 1, 0), (0, 0, 2)$	\longleftrightarrow	$R_1 = (10, 10, 02)$	\longleftrightarrow	$\mathbf{d}_1 = (0, 0, 0)$
$(1, 1, 0), (0, 1, 2)$	\longleftrightarrow	$R_2 = (10, 11, 02)$	\longleftrightarrow	$\mathbf{d}_2 = (0, 1, 0)$
$(0, 0, 1), (1, 0, 1)$	\longleftrightarrow	$R_3 = (01, 00, 11)$	\longleftrightarrow	$\mathbf{d}_3 = (1, 0, 0)$
$(1, 1, 0), (1, 0, 1)$	\longleftrightarrow	$R_4 = (11, 10, 01)$	\longleftrightarrow	$\mathbf{d}_4 = (2, 0, 0)$
$(0, 0, 1), (0, 0, 1)$	\longleftrightarrow	$R_5 = (00, 00, 11)$	\longleftrightarrow	$\mathbf{d}_5 = (3, 0, 0)$
$(1, 1, 0), (0, 0, 1)$	\longleftrightarrow	$R_6 = (10, 10, 01)$	\longleftrightarrow	$\mathbf{d}_6 = (0, 0, 1)$
$(0, 0, 1), (2, 0, 0)$	\longleftrightarrow	$R_7 = (02, 00, 10)$	\longleftrightarrow	$\mathbf{d}_7 = (4, 0, 0)$
$(0, 0, 0), (1, 0, 0)$	\longleftrightarrow	$R_8 = (01, 00, 00)$	\longleftrightarrow	$\mathbf{d}_8 = (1, 0, 1)$
$(1, 1, 0), (0, 0, 0)$	\longleftrightarrow	$R_9 = (10, 10, 00)$	\longleftrightarrow	$\mathbf{d}_9 = (0, 0, 2)$
$(0, 0, 0), (0, 1, 0)$	\longleftrightarrow	$R_{10} = (00, 01, 00)$	\longleftrightarrow	$\mathbf{d}_{10} = (1, 1, 0)$
$(0, 0, 0), (0, 0, 0)$	\longleftrightarrow	$R_{11} = (00, 00, 00)$	\longleftrightarrow	$\mathbf{d}_{11} = (2, 0, 1)$
$(0, 0, 0), (0, 0, 1)$	\longleftrightarrow	$R_{12} = (00, 00, 01)$	\longleftrightarrow	$\mathbf{d}_{12} = (1, 0, 2)$
$(1, 1, 0), (0, 1, 0)$	\longleftrightarrow	$R_{13} = (10, 11, 00)$	\longleftrightarrow	$\mathbf{d}_{13} = (0, 1, 1)$
$(1, 1, 0), (1, 0, 0)$	\longleftrightarrow	$R_{14} = (11, 10, 00)$	\longleftrightarrow	$\mathbf{d}_{14} = (3, 0, 1)$
$(0, 0, 1), (0, 0, 0)$	\longleftrightarrow	$R_{15} = (00, 00, 10)$	\longleftrightarrow	$\mathbf{d}_{15} = (0, 0, 3)$
$(1, 1, 0), (0, 1, 1)$	\longleftrightarrow	$R_{16} = (10, 11, 01)$	\longleftrightarrow	$\mathbf{d}_{16} = (0, 1, 2)$
$(1, 1, 0), (1, 1, 0)$	\longleftrightarrow	$R_{17} = (11, 11, 00)$	\longleftrightarrow	$\mathbf{d}_{17} = (2, 1, 0)$
$(0, 0, 1), (1, 0, 0)$	\longleftrightarrow	$R_{18} = (01, 00, 10)$	\longleftrightarrow	$\mathbf{d}_{18} = (2, 0, 2)$
$(1, 1, 1), (0, 1, 0)$	\longleftrightarrow	$R_{19} = (10, 11, 10)$	\longleftrightarrow	$\mathbf{d}_{19} = (1, 0, 3)$
$(1, 1, 1), (0, 0, 0)$	\longleftrightarrow	$R_{20} = (10, 10, 10)$	\longleftrightarrow	$\mathbf{d}_{20} = (0, 1, 3)$

Therefore, a monomial linear basis of $K[X]/\mathfrak{S}(\wp)$ consists of the equivalence classes (modulo $\mathfrak{S}(\wp)$) of the monomials

1	x_1	x_1^2	x_1^3	x_1^4	x_3	x_1x_3	$x_1^2x_3$	$x_1^3x_3$
x_2	x_1x_2	$x_1^2x_2$			x_2x_3			
x_3^2	$x_1x_3^2$	$x_1^2x_3^2$			x_3^3	$x_1x_3^3$		
$x_2x_3^2$					$x_2x_3^3$			

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