

Finding All q -Hypergeometric Solutions of q -Difference Equations

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Abstract

We present an algorithm for finding all solutions $y(x)$ of a linear homogeneous q -difference equation such that $y(qx)/y(x)$ is a rational function of q and x . The algorithm can also be used to construct q -hypergeometric series solutions of q -difference equations.

Résumé

Nous présentons un algorithme qui trouve toutes les solutions $y(x)$ des équations linéaires homogènes aux q -différences, telles que $y(qx)/y(x)$ est une fonction rationnelle de q et de x . On peut utiliser cet algorithme aussi pour construire les solutions des équations aux q -différences ayant la forme d'une série q -hypergéométrique.

1 Introduction

Let \mathbb{Q} be the rational number field, q transcendental over \mathbb{Q} , K a computable extension of $\mathbb{Q}(q)$, and x transcendental over K . Denote by Q the unique automorphism of $K(x)$ which fixes K and satisfies $Qx = qx$. Then $K(x)$ together with Q is an inversive difference field.

Let M be a difference extension ring of $K(x)$. An element $a \in M$ is q -polynomial if $a \in K[x]$, and q -rational if $a \in K(x)$. An element $a \in M \setminus \{0\}$ is a q -hypergeometric term if $Qa = ra$ for some $r \in K(x)$. All these concepts are relative to the field K .

We are interested in q -hypergeometric solutions y of $Ly = 0$ where

$$L = \sum_{i=0}^{\rho} p_i Q^i$$

is a linear q -difference operator of order ρ with coefficients $p_i \in K(x)$, with $p_\rho, p_0 \neq 0$. By clearing denominators in $Ly = 0$ we can restrict our attention to operators L with $p_i \in K[x]$. An algorithm for this problem is presented in Section 4. It is a q -analogue of the algorithm

for finding hypergeometric solutions of difference equations described in [6]. In preparation, we show how to find q -polynomial solutions of $Ly = 0$ in Section 2, and give a normal form for q -rational functions in Section 3. Finally, in Section 5, we describe solution of various related problems such as solving nonhomogeneous equations, finding solutions in the form of q -hypergeometric series, and deriving q -hypergeometric identities.

We use \mathbb{N} to denote the set of nonnegative integers. By $(a; q)_n$ we denote the expression $(1 - a)(1 - aq) \cdots (1 - aq^{n-1})$.

In our examples we use two algebraic settings which are special cases of the general framework described above. In one we work with sequences of elements of K , identifying sequences which agree from some point on. More precisely, we take $M = K^{\mathbb{N}}/J$ where $K^{\mathbb{N}}$ is the ring of sequences over K , and J is the ideal of sequences with finitely many nonzero terms. In particular, all equalities among sequences (of the form $a_n = b_n$) are meant to hold for all but finitely many $n \in \mathbb{N}$. Further we take $x = (q^n)_{n=0}^{\infty} + J$ and define Q as the unique automorphism of M satisfying $Q(a + J) = Ea + J$ for all $a \in K^{\mathbb{N}}$. Here E denotes the shift operator acting on $K^{\mathbb{N}}$ by $Ea_n = a_{n+1}$. Obviously K can be embedded in M as the subring of constant sequences. To simplify notation, we will henceforth identify $a + J \in K^{\mathbb{N}}/J$ with its representative $a \in K^{\mathbb{N}}$. Note that in this context a sequence a_n is q -polynomial if $a_n = p(q^n)$ for some $p \in K[x]$, q -rational if $a_n = r(q^n)$ for some $r \in K(x)$, and a q -hypergeometric term if $a_{n+1} = r(q^n)a_n$ for some $r \in K(x)$.

In another setting we take $M = K[[x]]$ (or $M = K((x))$), the ring of formal power series (resp. the field of formal Laurent series) over K . Again, K , $K[x]$, and $K(x)$ are embedded in M in a natural way. We distinguish between series that are q -hypergeometric terms, and series whose coefficients form a q -hypergeometric sequence. More precisely, a series $f(x) = \sum_{j=0}^{\infty} \alpha_j x^j$ is a q -hypergeometric term if $f(qx) = r(x)f(x)$ for some $r(x) \in K(x)$, and a q -hypergeometric series if $\alpha_{j+1} = r(q^j)\alpha_j$ for some $r(x) \in K(x)$ and for all large enough $j \in \mathbb{N}$.

Several times we will need to find the largest $n \in \mathbb{N}$ (if any) such that q^n is a root of a given polynomial with coefficients in K . Therefore we assume that K is a q -suitable field, meaning that there exists an algorithm which given $p \in K[x]$ finds all $n \in \mathbb{N}$ such that $p(q^n) = 0$. For instance, if $K = k(q)$ where q is transcendental over k we can proceed as follows: Let $p(x) = \sum_{i=0}^d c_i x^i$ where $c_i \in k[q]$. Compute $s = \min\{i; c_i \neq 0\}$ and $t = \max\{j; q^j | c_s\}$. Then $p(q^n) = 0$ only if $n \leq t$, and the set of all such n can be found by testing the values $n = t, t - 1, \dots, 0$.

2 q -polynomial solutions

First we show how to find solutions $y \in K[x]$ of $Ly = 0$. Let $p_i = \sum_{k=0}^d c_{ik} x^k$ where $c_{ik} \in k[q]$ and not all c_{id} are zero. Assume that $y = \sum_{j=0}^N \alpha_j x^j$ where $\alpha_N \neq 0$. Substituting these expressions into $Ly = 0$ and replacing k by $l = j + k$ yields

$$\sum_{i,l,j} c_{i,l-j} \alpha_j q^{ij} x^l = 0$$

which implies that

$$\sum_{j=\max\{l-d,0\}}^{\min\{l,N\}} \sum_{i=0}^{\rho} c_{i,l-j} \alpha_j q^{ij} = 0, \quad \text{for } 0 \leq l \leq N + d. \quad (1)$$

In particular, for $l = N + d$,

$$\sum_{i=0}^{\rho} c_{id} q^{iN} = 0, \quad (2)$$

and for $l = 0$,

$$\alpha_0 \sum_{i=0}^{\rho} c_{i0} = 0. \quad (3)$$

From (2) it follows that q^N is a root of the polynomial $P(x) = \sum_{i=0}^{\rho} c_{id} x^i$. Let N_0 be the largest $n \in \mathbb{N}$ such that $P(q^n) = 0$ (see the last paragraph of Introduction). All q -polynomial solutions y of $Ly = 0$ can now be found by the method of undetermined coefficients. Ultimately, the problem is reduced to a system of linear algebraic equations over K with $N_0 + 1$ unknowns. – A more efficient method leading to a system with at most $\min\{2d, N_0 + 1\}$ unknowns is described in [2].

3 A normal form for q -rational functions

Theorem 1 *Let $r \in K(x) \setminus \{0\}$. Then there are $z \in K$ and monic polynomials $a, b, c \in K[x]$ such that*

$$r(x) = z \frac{a(x)}{b(x)} \frac{c(qx)}{c(x)}, \quad (4)$$

$$\gcd(a(x), b(q^n x)) = 1 \quad \text{for all } n \in \mathbb{N}, \quad (5)$$

$$\gcd(a(x), c(x)) = 1, \quad (6)$$

$$\gcd(b(x), c(qx)) = 1, \quad (7)$$

$$c(0) \neq 0. \quad (8)$$

Proof: Write $r(x) = \frac{f(x)}{g(x)}$ where f, g are relatively prime polynomials. We start by finding the set \mathcal{S} of all $n \in \mathbb{N}$ such that $f(x)$ and $g(q^n x)$ have a nonconstant common factor. To this end consider the polynomial $R(h) = \text{Resultant}_x(f(x), g(hx))$. By the well-known properties of polynomial resultants, $\mathcal{S} = \{n \in \mathbb{N}; R(q^n) = 0\}$.

Assume that $\mathcal{S} = \{n_1, n_2, \dots, n_t\}$ where $t \geq 0$ and $n_1 < n_2 < \dots < n_t$. In addition, let $n_{t+1} = +\infty$. Define polynomials f_i and g_i inductively by setting

$$f_0(x) = f(x), \quad g_0(x) = g(x),$$

and for $i = 1, 2, \dots, t$,

$$s_i(x) = \gcd(f_{i-1}(x), g_{i-1}(q^{n_i} x)),$$

$$f_i(x) = f_{i-1}(x) / s_i(x),$$

$$g_i(x) = g_{i-1}(x) / s_i(q^{-n_i} x).$$

Now take

$$\begin{aligned} z &= \alpha/\beta, \\ a(x) &= f_t(x)/\alpha, \\ b(x) &= g_t(x)/\beta, \\ c(x) &= \prod_{i=1}^t \prod_{j=1}^{n_i} s_i(q^{-j}x), \end{aligned}$$

where α and β denote the leading coefficients of $f_t(x)$ and $g_t(x)$, respectively. Before proving (4) – (8) we state a lemma.

Lemma 1 *Let $n \in \mathbb{N}$. If $0 \leq l \leq i, j \leq t$ and $n < n_{l+1}$, then $\gcd(f_i(x), g_j(q^n x)) = 1$.*

Proof: Assume first that $n \notin \mathcal{S}$. Then $R(q^n) \neq 0$, hence $\gcd(f(x), g(q^n x)) = 1$. Since $f_i(x) | f(x)$ and $g_j(x) | g(x)$ it follows that $\gcd(f_i(x), g_j(q^n x)) = 1$, too.

To prove the lemma for $n \in \mathcal{S}$ we use induction on l .

$l = 0$: In this case there is nothing to prove since there is no $n \in \mathcal{S}$ such that $n < n_1$.

$l > 0$: Assume that the lemma holds for all $n < n_l$. It remains to show that it also holds for $n = n_l$. Since $f_i(x) | f_l(x)$ and $g_j(x) | g_l(x)$ it follows that $\gcd(f_i(x), g_j(q^{n_l} x))$ divides $\gcd(f_l(x), g_l(q^{n_l} x)) = \gcd(f_{l-1}(x)/s_l(x), g_{l-1}(q^{n_l} x)/s_l(x))$. By the definition of $s_l(x)$, the latter gcd is 1, completing the proof. \square

Now we proceed to verify properties (4) – (8).

(4):

$$\begin{aligned} z \frac{a(x) c(qx)}{b(x) c(x)} &= \frac{f_t(x)}{g_t(x)} \prod_{i=1}^t \prod_{j=1}^{n_i} \frac{s_i(q^{1-j}x)}{s_i(q^{-j}x)} \\ &= \frac{f_0(x)}{\prod_{i=1}^t s_i(x)} \frac{\prod_{i=1}^t s_i(q^{-n_i}x)}{g_0(x)} \prod_{i=1}^t \frac{s_i(x)}{s_i(q^{-n_i}x)} = \frac{f(x)}{g(x)} = r(x). \end{aligned}$$

(5): Let $i = j = l = t$ in Lemma 1. Then $\gcd(f_t(x), g_t(q^n x)) = 1$ for all $n < n_{t+1} = +\infty$. In other words, $\gcd(a(x), b(q^n x)) = 1$ for all $n \in \mathbb{N}$.

(6): If $a(x)$ and $c(x)$ have a non-constant common factor then so do $f_t(x)$ and $s_i(q^{-j}x)$, for some i and j such that $1 \leq i \leq t$ and $1 \leq j \leq n_i$. Since $g_{i-1}(q^{n_i-j}x) = g_i(q^{n_i-j}x)s_i(q^{-j}x)$, it follows that $g_{i-1}(q^{n_i-j}x)$ contains this factor as well. As $n_i - j < n_i$, this contradicts Lemma 1. Hence $a(x)$ and $c(x)$ are relatively prime.

(7): If $b(x)$ and $c(qx)$ have a non-constant common factor then so do $g_t(x)$ and $s_i(q^{-j}x)$, for some i and j such that $1 \leq i \leq t$ and $1 \leq j+1 \leq n_i$. Since $f_{i-1}(q^{-j}x) = f_i(q^{-j}x)s_i(q^{-j}x)$, it follows that $f_{i-1}(x)$ and $g_t(q^j x)$ contain this factor as well. As $j < n_i$, this contradicts Lemma 1. Hence $b(x)$ and $c(qx)$ are relatively prime.

(8): It is easy to see that $s_i(x)$ divides both $f(x)$ and $g(q^{n_i}x)$. Hence $s_i(0) = 0$ would imply that $f(0) = g(0) = 0$, contrary to the assumption that f and g are relatively prime. It follows that $s_i(0) \neq 0$ for all i , and consequently $c(0) \neq 0$. \square

Theorem 2 Let $a, b, c, A, B, C \in K[x]$ be polynomials such that $c(0) \neq 0$ and $\gcd(a(x), c(x)) = \gcd(b(x), c(qx)) = \gcd(A(x), B(q^n x)) = 1$, for all $n \in \mathbb{N}$. If

$$\frac{a(x) c(qx)}{b(x) c(x)} = \frac{A(x) C(qx)}{B(x) C(x)}, \quad (9)$$

then $c(x)$ divides $C(x)$.

Proof: Let

$$\begin{aligned} g(x) &= \gcd(c(x), C(x)), \\ d(x) &= c(x)/g(x), \\ D(x) &= C(x)/g(x). \end{aligned}$$

Then $\gcd(d(x), D(x)) = \gcd(a(x), d(x)) = \gcd(b(x), d(qx)) = 1$ and $d(0) \neq 0$. Clear denominators in (9) and cancel $g(x)g(qx)$ on both sides. The result $A(x)b(x)d(x)D(qx) = a(x)B(x)D(x)d(qx)$ shows that

$$\begin{aligned} d(x) &| B(x)d(qx), \\ d(qx) &| A(x)d(x). \end{aligned}$$

Using these two relations repeatedly we find that

$$\begin{aligned} d(x) &| B(x)B(qx) \cdots B(q^{n-1}x)d(q^n x), \\ d(x) &| A(q^{-1}x)A(q^{-2}x) \cdots A(q^{-n}x)d(q^{-n}x), \end{aligned}$$

for all $n \in \mathbb{N}$. It is easy to see that since $d(0) \neq 0$ and q is not a root of unity, $d(x)$ and $d(q^n x)$ are relatively prime for all large enough n . It follows that $d(x)$ divides both $B(x)B(qx) \cdots B(q^{n-1}x)$ and $A(q^{-1}x)A(q^{-2}x) \cdots A(q^{-n}x)$ for all large enough n . But these polynomials are relatively prime by assumption, so $d(x)$ is constant. Hence $c(x) | C(x)$. \square

Corollary 1 The factorization of $r(x)$ described in Theorem 1 is unique.

Proof: If

$$z \frac{a(x) c(qx)}{b(x) c(x)} = Z \frac{A(x) C(qx)}{B(x) C(x)}$$

are two such factorizations then $c(x) | C(x)$ and $C(x) | c(x)$, by Theorem 2. Since these polynomials are monic, $c = C$. It follows that $z = Z$ and $aB = Ab$. Hence $a | A$ and $A | a$, so $a = A$ and $b = B$. \square

Corollary 2 Among all factorizations of $r(x)$ satisfying (4) and (5) of Theorem 1, the one satisfying (4) - (8) has $c(x)$ of least degree.

4 q -hypergeometric solutions

After this preparation we turn to the algorithm for finding q -hypergeometric solutions y of $Ly = 0$. Let $Qy = ry$ where $r \in K(x)$, then $Q^i y = \prod_{j=0}^{i-1} r(q^j x)y$. We look for $r(x)$ in the normal form described in Theorem 1. After inserting (4) into $Ly = 0$, clearing denominators and cancelling y we obtain

$$\sum_{i=0}^{\rho} z^i f_i(x) c(q^i x) = 0 \quad (10)$$

where

$$f_i(x) = p_i(x) \prod_{j=0}^{i-1} a(q^j x) \prod_{j=i}^{\rho-1} b(q^j x).$$

Since all terms in (10) except for $i = 0$ are divisible by $a(x)$ it follows that $a(x)$ divides $p_0(x) \prod_{j=0}^{\rho-1} b(q^j x) c(x)$. Because of (5) and (6), $a(x)$ divides $p_0(x)$. Similarly, all terms in (10) except for $i = \rho$ are divisible by $b(q^{\rho-1} x)$, therefore $b(q^{\rho-1} x)$ divides $z^{\rho} p_{\rho}(x) \prod_{j=0}^{\rho-1} a(q^j x) c(q^{\rho} x)$. Because of (5) and (7), $b(q^{\rho-1} x)$ divides $p_{\rho}(x)$. Thus we have a finite choice for $a(x)$ and $b(x)$.

For each choice of $a(x)$ and $b(x)$, equation (10) is a q -difference equation for the unknown polynomial $c(x)$. However, $z \in K$ is also not known yet. Let u_{ik} denote the coefficient of x^k in f_i . Since $c(0) \neq 0$, we have $\alpha_0 \neq 0$ in (3), hence applying (3) to (10) we obtain

$$\sum_{i=0}^{\rho} u_{i0} z^i = 0. \quad (11)$$

We may assume that not all u_{i0} are zero, or else we start by first cancelling a power of x from the coefficients of (10). Thus z is a nonzero root of $f(z) = \sum_{i=0}^{\rho} u_{i0} z^i$, and is algebraic over K .

If $N = \deg c(x)$ then by (2),

$$\sum_{i=0}^{\rho} u_{id} z^i q^{iN} = 0, \quad (12)$$

hence $w = zq^N$ is a nonzero root of $g(w) = \sum_{i=0}^{\rho} u_{id} w^i$. It follows that q^N is a root of $p(x) = \text{Resultant}_w(f(w), g(wx))$, thus to obtain an upper bound on N computation in algebraic extensions of K is not necessary.

In summary, we find the factors of $r(x)$ as follows:

1. $a(x)$ is a monic factor of $p_0(x)$,
2. $b(x)$ is a monic factor of $p_{\rho}(q^{1-\rho} x)$,
3. z is a root of Eqn. (11),
4. $c(x)$ is a nonzero q -polynomial solution of (10).

Then $r = z(a/b)(Qc/c)$ and $Qy = ry$.

Example 1 Let us find a q -hypergeometric solution y of $Ly = 0$ where

$$L = xQ^3 - q^3x^2Q^2 - (x^2 + q)Q + qx(x^2 + q).$$

The candidates for $a(x)$ are

$$1, x, x^2 + q, x(x^2 + q),$$

and the candidates for $b(x)$ are

$$1, x.$$

Here we explore only the choice $a(x) = x$ and $b(x) = 1$. The corresponding equation (10) is, after cancelling one x ,

$$z^3q^3x^3c(q^3x) - z^2q^4x^3c(q^2x) - z(x^2 + q)c(qx) + q(x^2 + q)c(x) = 0, \quad (13)$$

whence $f(z) = -qz + q^2$ with unique root $z = q$, and $g(w) = q^3w^3 - q^4w^2$ with unique nonzero root $w = q = zq^N = q^{N+1}$. It follows that $N = 0$ is the only possible degree for c . Equation (13) is satisfied by $c = 1$. Thus we have found $r = z(a/b)(Qc/c) = qx$, and the corresponding q -hypergeometric solution of $Ly = 0$ satisfies $Qy = qxy$. We can take, for instance, $y_n = x(x/q)(x/q^2) \cdots (x/q^n) = q^{\binom{n+1}{2}}$.

To find other q -hypergeometric solutions (if any), the remaining combinations for $a(x)$ and $b(x)$ could be tried; or even better, the order of the equation could be reduced using the obtained solution, and the algorithm used recursively on the reduced equation. Our *Mathematica* implementation of this algorithm (which we call `qHyper`) shows that up to a constant factor, there are in fact no other q -hypergeometric solutions:

```
In[1]:= qHyper[x y[q^3 x] - q^3 x^2 y[q^2 x] -
          (x^2 + q) y[q x] + q x (x^2 + q) y[x] == 0, y[x]]
Out[1]= {q x}
```

Note that `qHyper` returns a list of quotients Qy/y rather than solutions y themselves. \square

Example 2 Consider the equation $Ly = 0$ where $L = Q^2 - (1 + q)Q + q(1 - qx^2)$. As shown by `qHyper`,

```
In[2]:= qHyper[y[q^2 x] - (1 + q) y[q x] + q (1 - q x^2) y[x], y[x]]
Out[2]= {1 - Sqrt[q] x, 1 + Sqrt[q] x}
```

this equation has two linearly independent q -hypergeometric solutions, $(\sqrt{q}; q)_n$ and $(-\sqrt{q}; q)_n$. Here K is the splitting field of $1 - qx^2$. \square

5 Some related problems

5.1 Nonhomogeneous equations

Consider the problem of finding q -hypergeometric solutions y of the nonhomogeneous equation $Ly = b$ where $b \neq 0$. Let $Qy = ry$ where $r \in K(x)$. Then $Ly = fy$ where $f = \sum_{i=0}^p p_i \prod_{j=0}^{i-1} Q^j r \in K(x)$. This simple fact has two important consequences:

1. $b = fy$ is q -hypergeometric,
2. $y = b/f$ is a q -rational multiple of b .

Let $Qb = sb$ where $s \in K(x)$ is given. We look for y in the form $y = fb$ where $f \in K(x)$ is an unknown q -rational function. Substituting this into $Ly = b$ gives

$$\sum_{i=0}^{\rho} p_i \left(\prod_{j=0}^{i-1} Q^j s \right) Q^i f = 1.$$

Now q -rational solutions of this equation can be found using the algorithm given in [1].

In particular, this gives an algorithm for the problem of *indefinite q -hypergeometric summation*: Given a q -hypergeometric sequence b_n , decide if $y_n = \sum_{j=0}^{n-1} b_j$ is q -hypergeometric, and if so, express it in closed form. Obviously y_n satisfies $y_{n+1} - y_n = b_n$. Since we are interested in q -hypergeometric solutions, we can rewrite this as $Qy - y = b$ and use the technique described above.

Example 3 Let $y_n = \sum_{j=0}^{n-1} b_j$ where $b_n = q^n(q; q)_n$. Then y satisfies the equation

$$Qy - y = b \tag{14}$$

where $s = Qb/b = q(1 - qx)$. The equation for f is

$$q(1 - qx)Qf - f = 1,$$

with unique q -rational solution $f = -1/(qx)$. Hence $y_n = C - (q; q)_n/q$ where C is a constant. Since $y_0 = 0$ it follows that $C = 1/q$ and $y_n = (1 - (q; q)_n)/q$. \square

The same technique for solving nonhomogeneous equations also works when we look for q -hypergeometric term solutions in $M = K[[x]]$.

Example 4 Let

$$Q^2y(x) - (1 - qx)Qy(x) + qy(x) = b(x) \tag{15}$$

where

$$b(x) = \sum_{i=0}^{\infty} \frac{x^i}{(q; q)_i}.$$

Here $b(qx) = (1 - x)b(x)$, as can be easily verified. Thus $s = 1 - x$ and the equation for f is

$$(1 - qx)(1 - x)Q^2f - (1 - qx)(1 - x)Qf + qf = 1$$

with q -rational solution $f = 1/q$. Hence $y(x) = b(x)/q$ solves (15). \square

5.2 q -hypergeometric series solutions

Assume that $y = \sum_{j=0}^{\infty} \alpha_j x^j$ and $Ly = b$ where $b = \sum_{j=0}^{\infty} \beta_j x^j$. As in (1), we obtain

$$\sum_{j=\max\{l-d,0\}}^l \sum_{i=0}^{\rho} c_{i,l-j} \alpha_j q^{ij} = \beta_l, \quad \text{for } l \geq 0. \quad (16)$$

We separate the cases $0 \leq l < d$ and $l \geq d$. In the former case, (16) yields initial conditions

$$\sum_{j=0}^l \alpha_j \sum_{i=0}^{\rho} c_{i,l-j} q^{ij} = \beta_l, \quad \text{for } 0 \leq l < d, \quad (17)$$

while in the latter, substitutions $m = l - d$, $s = j - m$, and $X = q^m$ transform (16) into the associated q -difference equation

$$\sum_{s=0}^d \alpha_{m+s} \sum_{i=0}^{\rho} c_{i,d-s} q^{is} X^i = \beta_{m+d}, \quad \text{for } m \geq 0, \quad (18)$$

for the unknown sequence $(\alpha_m)_{m=0}^{\infty}$. We use the algorithms of Sections 4 and 5.1 to find all solutions of (18) which are linear combinations of q -hypergeometric terms, then select the constants in these combinations so that conditions (17) are satisfied (if possible).

Example 5 Let us find q -hypergeometric series solutions y of

$$q^2 x^2 Q^3 y + (1 + q)xQ^2 y + (1 - x)Qy - y = 0. \quad (19)$$

The associated equation (18) in this case is

$$(q^2 X - 1)\alpha_{m+2} + (q^2(q+1)X^2 - qX)\alpha_{m+1} + q^2 X^3 \alpha_m = 0 \quad (20)$$

and `qHyper` finds two solutions:

```
In[3] := qHyper[(q^2 X - 1) y[q^2 X] + (q^2 (1 + q) X^2 - q X) y[q X] +
              q^2 X^3 y[X] == 0, y[X]]
```

```
Out[3] = {-X,  $\frac{q X^2}{1 - q X}$ }
```

Thus the general solution of (20) is $\alpha_m = Cq^{m^2}/(q; q)_m + D(-1)^m q^{\binom{m}{2}}$ where C and D are arbitrary constants. Equations (17) imply that $D = 0$. Hence $y^{(1)} = \sum_{m=0}^{\infty} q^{m^2} x^m / (q; q)_m$ is a q -hypergeometric series solution of (19).

Note that running `qHyper` on equation (19) itself we obtain another solution $y^{(2)} = (-1)^n / q^{\binom{n}{2}}$. \square

Example 6 The right-hand side of the equation (15) is both a q -hypergeometric term and a q -hypergeometric series. The associated nonhomogeneous equation

$$(qX^2 - X + 1)\alpha_{m+1} + X\alpha_m = \frac{1}{q(q; q)_{m+1}}$$

can be solved as described in Section 5.1. Here $s = 1/(1 - q^2X)$ and the equation for f

$$\frac{1 - X + qX^2}{1 - q^2X}Qf + Xf = 1$$

is satisfied by the q -rational function $f = 1 - qX$. Thus $\alpha_m = (1 - qX)/(q(q; q)_{m+1}) = 1/(q(q; q)_m)$, and we find the same solution $y(x) = b(x)/q$ as in Example 4. \square

5.3 Deriving q -hypergeometric identities

Another important application is definite q -hypergeometric summation. The corresponding algorithm of [7] will produce a q -difference equation for the sum, but in general it will not be of minimal order. Thus it can happen that the equation will be of order 2 or more while the sum can actually be expressed in closed form. In this case one can use our algorithm to find the q -hypergeometric solutions of the equation, and then test them to see which linear combination – if any – gives the initial sum.

In analogy with the ordinary hypergeometric case [4], we also expect our algorithm to play an important role in the factorization algorithm for linear q -difference operators.

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